

Periods of modular functions and diophantine approximation

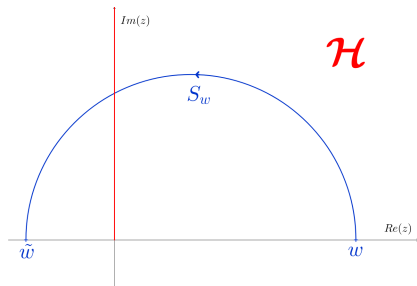
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April 14, 2020

Periods of modular functions

$$\Gamma = \mathrm{SL}(2, \mathbb{Z})$$



$$Q(x, y) = ax^2 + bxy + cy^2 \in \mathbb{Z}[x, y], \quad (a, b, c) = 1$$

$$b^2 - 4ac = D > 0, \quad D \neq \square$$

$$w = \frac{-b + \sqrt{D}}{2a}, \quad \tilde{w} = \frac{-b - \sqrt{D}}{2a}$$

$$S_w : a|z|^2 + b\mathrm{Re}(z) + c = 0 \quad (z \in \mathcal{H})$$

Stabilizer of Q in Γ : $\Gamma_w = \langle A_w \rangle$ where

$A_w = \begin{pmatrix} \frac{1}{2}(t - bu) & -cu \\ au & \frac{1}{2}(t + bu) \end{pmatrix}$, with (t, u) the smallest primitive solution to Pell's equation $t^2 - Du^2 = 4$.

$$\begin{aligned} \Gamma_w &\xrightarrow{\cong} \left\{ \text{units of norm 1 of } \mathbb{Q}(\sqrt{D}) \right\}^2 \\ A_w &\mapsto \varepsilon_D^2 = ((t + u\sqrt{D})/2)^2 \end{aligned}$$

$$j(w) := \frac{1}{\text{length}(S_w)} \int_{\Gamma_w \backslash S_w} j(z) \frac{\sqrt{D}}{Q(z, 1)} dz$$

$$\text{length}(S_w) = \int_{\Gamma_w \backslash S_w} \frac{\sqrt{D}}{Q(z, 1)} dz = 2 \log \varepsilon_D$$

Choose any point z_0 in \mathcal{H} ,

$$\int_{\Gamma_w \backslash S_w} = \int_{z_0}^{A_w z_0}$$

Distribution of periods of j

$$Tr_D(j) = \sum_{Q \in \mathcal{Q}_D/\Gamma} j(w_Q), \quad \mathcal{Q}_D = \{\text{bin. quad. forms of discr. } D\}$$

Duke-Friedlander-Iwaniec and independently Masri:

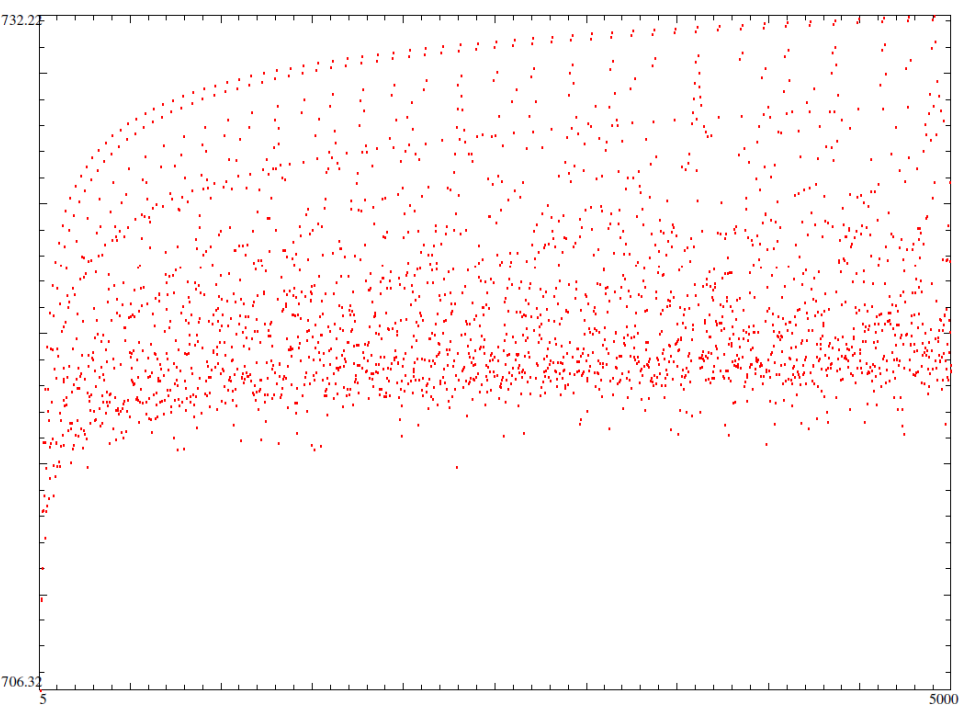
$$\frac{Tr_D(j)}{h^+(D)} \rightarrow 720 \text{ as } D \rightarrow +\infty.$$

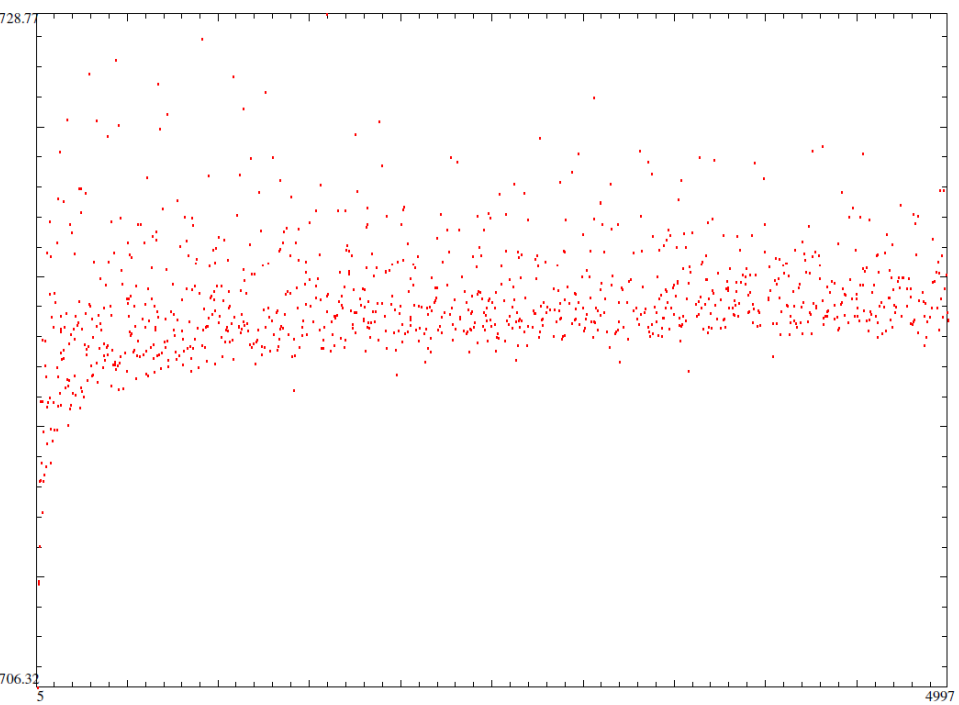
Theorem (B-Imamoglu)

[Conjectures of Kaneko]

- (i) $\text{Re}(j(w)) \leq 744$, $(j(z) = q^{-1} + 744 + \dots)$
and the bound is optimal;
- (ii) If all the partial quotients in the period of the (negative) continued fraction expansion of w are $\geq 3e^{55}$, then

$$\text{Re}(j(w)) \geq j\left(\frac{1 + \sqrt{5}}{2}\right) \approx 706.3$$





For $N > 2$, consider

$$\frac{N + \sqrt{N^2 - 4}}{2} = (N, N, N, \dots) = (\overline{N}),$$

where

$$(a_0, a_1, a_2, \dots) = a_0 - \frac{1}{a_1 - \frac{1}{a_2 - \frac{1}{\ddots}}}$$

with $(a_0 \in \mathbb{Z}, a_i \geq 2 \text{ for } i \geq 1)$.

Theorem (B-I)

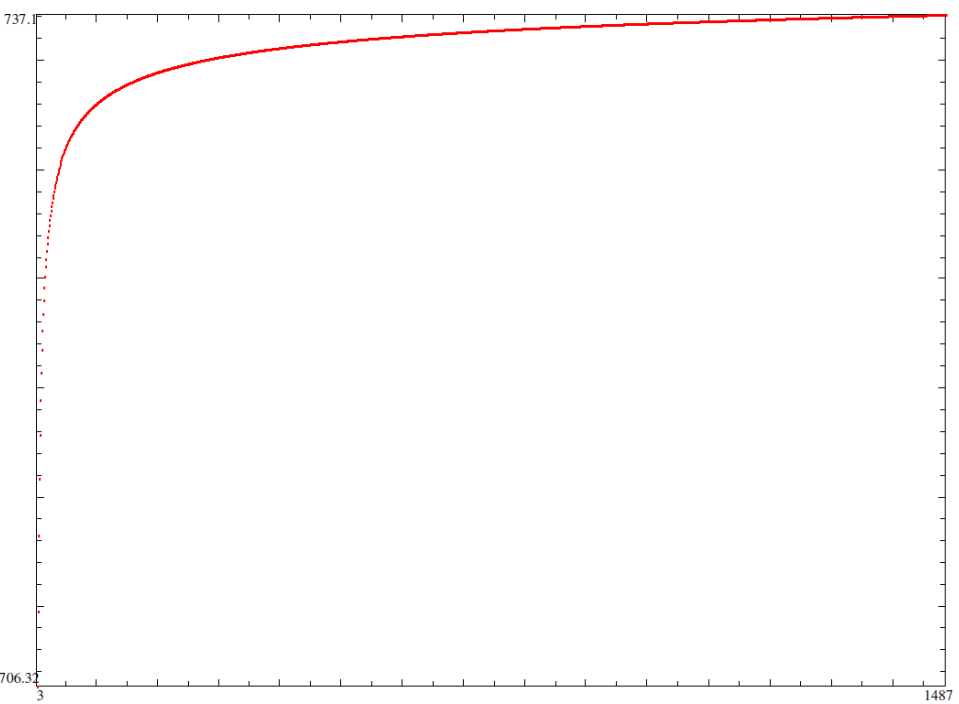
For all $N > 2$, $j\left(\frac{N + \sqrt{N^2 - 4}}{2}\right) \in \mathbb{R}$ and $\lim_{N \rightarrow \infty} j\left(\frac{N + \sqrt{N^2 - 4}}{2}\right) = 744$.

Theorem (B-I)

Let v and w be two quadratic irrationalities with respective periods in their c.f. $\overline{a_1, \dots, a_n}$ and $\overline{b_1, \dots, b_m}$ with $m|n$.

If $b_r \geq e^{55} a_r$ for all $r = 1, \dots, n$, then

$$\operatorname{Re}(j(w)) > \operatorname{Re}(j(v)).$$



Periods of j at Markov quadratics

Hurwitz: For all $x \in \mathbb{R}$,

$$\left| x - \frac{p}{q} \right| \leq \frac{1}{\sqrt{5}} \frac{1}{q^2}$$

for infinitely many $\frac{p}{q}$ with $(p, q) = 1$.

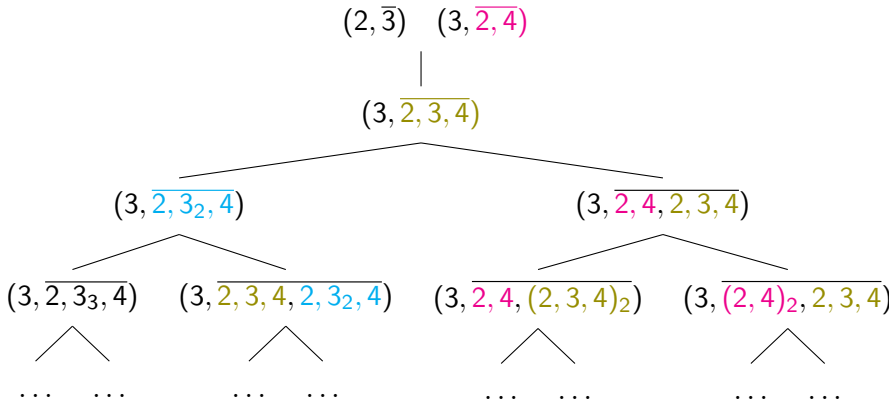
$$\nu(x) = \liminf_{q \rightarrow \infty} q \|qx\|, \quad L = \{\nu(x)\}_{x \in \mathbb{R}} \subseteq \left[0, \frac{1}{\sqrt{5}}\right]$$

- ▶ $\nu(x) = 0$ for almost all $x \in \mathbb{R}$
- ▶ if $\nu(x) > 0$, x is called **badly approximable**. $\nu(x) > 0 \Leftrightarrow$ all partial quotients of x are bounded. Example: **quadratic irrationalities** (their c.f. is eventually **periodic**)
- ▶ if x and x' are $\mathrm{PGL}(2, \mathbb{Z})$ -equivalent, then $\nu(x) = \nu(x')$

$L \cap [0, F]$ is continuous, where $F \approx 0.220856$ is Freiman's constant, $L \cap (F, \frac{1}{3}]$ has a fractal structure, $L \cap (\frac{1}{3}, \frac{1}{\sqrt{5}}]$ is discrete.

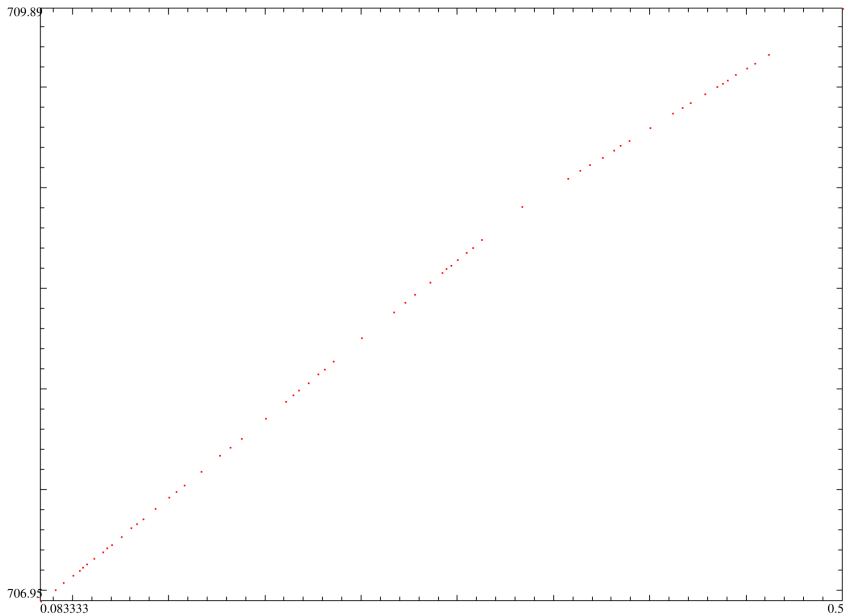
w is a Markov quadratic if $\nu(w) \in L \cap (\frac{1}{3}, \frac{1}{\sqrt{5}}]$ and any other x with $\nu(x) = \nu(w)$ is $\text{PGL}(2, \mathbb{Z})$ equivalent to w .

Markov quadratics:

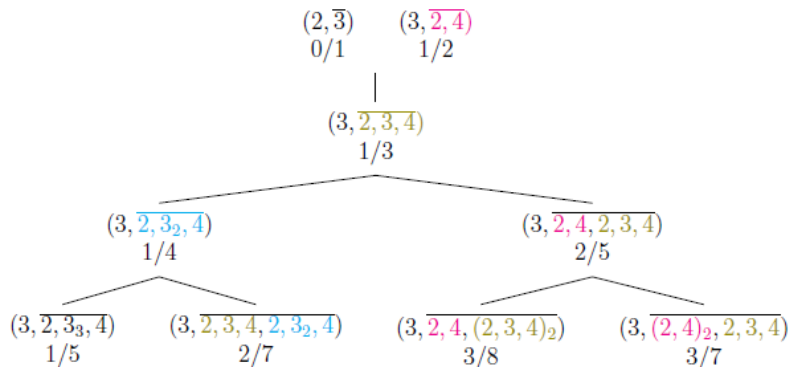


Theorem (B-I)

- (i) *The convergence property is preserved for the periods of j at the Markov quadratics.*
- (ii) *The sandwich property is preserved for the periods of j at the Markov quadratics that are located below some level in the tree.*



Farey parametrization of Markov tree



Why does Diophantine approximation come in?

Let w have a purely periodic c.f.

$$v_1 = w, \quad a_1 = \lceil v_1 \rceil; \quad v_{i+1} = \frac{1}{a_i - v_i}, \quad a_{i+1} = \lceil v_{i+1} \rceil$$

Then $w = (\overline{a_1, \dots, a_n})$ and $v_i = (\overline{a_i, a_{i+1}, \dots, a_n, \dots, a_{i-1}})$.

Another algorithm:

$$w_1 = w - 1, \quad w_{i+1} = \begin{cases} w_i - 1 = \begin{pmatrix} 1 & -1 \\ 0 & 1 \end{pmatrix} (w_i) & \text{if } w_i \geq 1, \\ \frac{w_i}{1 - w_i} = \begin{pmatrix} 1 & 0 \\ -1 & 1 \end{pmatrix} (w_i) & \text{if } 0 < w_i < 1. \end{cases}$$

This algorithm is **cyclic**: we get w_1, \dots, w_ℓ .

Lemma

$$j(w) = \frac{1}{2 \log \varepsilon_D} \int_{e^{\pi i/3}}^{e^{2\pi i/3}} j(z) K(z, w) dz,$$

where

$$K(z, w) = \sum_{i=1}^{\ell} \frac{1}{z - w_i} - \frac{1}{z - \tilde{w}_i}$$

and \tilde{w}_i is the Galois conjugate of w_i .

Example: proof of Th. 2.

$$w = (\overline{N}), \quad D = N^2 - 4, \quad \varepsilon_D = \frac{N + \sqrt{N^2 - 4}}{2}$$

$$w-1 = (N-1, \overline{N}) \xrightarrow{T^{-1}} (N-2, \overline{N}) \xrightarrow{T^{-1}} \dots \xrightarrow{T^{-1}} (1, \overline{N}) \xrightarrow{V^{-1}} (N-1, \overline{N})$$

$$K(z, w) = K_N(z) = \sum_{k=1}^{N-1} \frac{1}{z - (k, \overline{N})} - \frac{1}{z + (k, \overline{N})}$$

For $\theta \in [\pi/3, 2\pi/3]$,

- ▶ $\overline{K_N(e^{i\theta})} = K_N(e^{i(\pi-\theta)}) \Rightarrow \text{Im}(j(w)) = 0,$
- ▶ $K_N(e^{i\theta}) \rightarrow 2 \log N$ as $N \rightarrow \infty \Rightarrow \frac{K_N(e^{i\theta})}{2 \log \varepsilon_D} \rightarrow 1$ as $N \rightarrow \infty \Rightarrow$

$$\lim_{N \rightarrow \infty} j((\overline{N})) = \int_{\pi/3}^{2\pi/3} j(e^{i\theta}) i e^{i\theta} d\theta = 744.$$