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On the no-counterexample interpretation

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Abstract

In [15],[16] Kreisel introduced the no-counterexample interpretation (n.c.i.) of Peano arithmetic. In particular he proved, using a complicated ε -substitution method (due to W. Ackermann), that for every theorem A (A prenex) of first-order Peano arithmetic **PA** one can find ordinal recursive functionals Φ_A of order type $< \varepsilon_0$ which realize the Herbrand normal form A^H of A .

Subsequently more perspicuous proofs of this fact via functional interpretation (combined with normalization) and cut-elimination were found. These proofs however do not carry out the n.c.i. as a **local** proof interpretation and don't respect the modus ponens on the level of the n.c.i. of formulas A and $A \rightarrow B$. Closely related to this phenomenon is the fact that both proofs do not establish the condition (δ) and – at least not constructively – (γ) which are part of the definition of an ‘interpretation of a formal system’ as formulated in [15].

In this paper we determine the complexity of the n.c.i. of the modus ponens rule for

- (i) **PA**-provable sentences,
- (ii) for arbitrary sentences $A, B \in \mathcal{L}(\mathbf{PA})$ **uniformly** in functionals satisfying the n.c.i. of (prenex normal forms of) A and $A \rightarrow B$, and
- (iii) for arbitrary $A, B \in \mathcal{L}(\mathbf{PA})$ **pointwise** in given $\alpha (< \varepsilon_0)$ -recursive functionals satisfying the n.c.i. of A and $A \rightarrow B$.

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This yields in particular perspicuous proofs of new uniform versions of the conditions (γ) , (δ) .

Finally we discuss a variant of the concept of an interpretation presented in [17] and show that it is incomparable with the concept studied in [15],[16]. In particular we show that the n.c.i. of \mathbf{PA}_n by $\alpha(< \omega_n(\omega))$ -recursive functionals ($n \geq 1$) is an interpretation in the sense of [17] but not in the sense of [15] since it violates the condition (δ) .

1 Introduction

Let $\exists x A_0(x, \underline{a})$ be a Σ_1^0 -formula in the language $\mathcal{L}(\mathbf{PL})$ of first-order predicate logic \mathbf{PL} ($\underline{a} = a_1, \dots, a_k$ are all its free variables).

If

$$\mathbf{PL} \vdash \exists x A_0(x, \underline{a})$$

then by Herbrand's theorem there are terms $t_1[\underline{a}], \dots, t_n[\underline{a}]$ (built up out of \underline{a} , a distinguished object constant 0 and the object and function constants of A_0)¹ such that $\bigvee_{i=1}^n A_0(t_i[\underline{a}], \underline{a})$ is a tautology.

This extends to Σ_n^0 -formulas by introducing so-called index functions. For notational simplicity lets consider $n = 4$ only

$$A(\underline{a}) \equiv \exists x_1 \forall y_1 \exists x_2 \forall y_2 A_0(x_1, y_1, x_2, y_2, \underline{a}).$$

We replace y_1, y_2 by $f x_1, g x_1 x_2$, where f, g are new function symbols. If $\mathbf{PL} \vdash A$ then

$$\mathbf{PL} \vdash \exists x_1, x_2 A_0(x_1, f x_1, x_2, g x_1 x_2, \underline{a})$$

and so by Herbrand's theorem for Σ_1^0 -formulas there are terms built up from $\underline{a}, f, g, 0$ and the constants of $A_0(x, \underline{a})$ such that

$$\bigvee_{i=1}^n \bigvee_{j=1}^k A_0(t_i[\underline{a}, f, g], f(t_i[\underline{a}, f, g]), s_j[\underline{a}, f, g], g(t_i[\underline{a}, f, g], s_j[\underline{a}, f, g]), \underline{a})$$

is a tautology.

If we allow definition by cases and characteristic functions for quantifier-free formulas we can avoid the disjunction:

$$\Phi_1 \underline{a} f g := \begin{cases} t_1, & \text{if } \bigvee_{j=1}^k A_0(t_1, f(t_1), s_j, g(t_1, s_j), \underline{a}) \\ t_2 & \text{if } \neg (\text{case 1}) \wedge \bigvee_{j=1}^k A_0(t_2, \dots) \\ \vdots & \end{cases}$$

¹Throughout this paper A_0, B_0, C_0, \dots denote quantifier-free formulas.

$$\Phi_2 \underline{a} f g := \begin{cases} s_1, & \text{if } s_1 \text{ if } A_0(\Phi_1 \underline{a} f g, f(\Phi_1 \underline{a} f g), s_1, g(\Phi_1 \underline{a} f g, s_1), \underline{a}) \\ s_2 & \text{if } \neg (\text{case 1}) \wedge A_0(\Phi_1 \underline{a} f g, f(\Phi_1 \underline{a} f g), s_2, g(\Phi_1 \underline{a} f g, s_2), \underline{a}) \\ \vdots & \end{cases}$$

Then

$$(+)\ \forall \underline{a}, f, g\ A_0(\Phi_1 \underline{a} f g, f(\Phi_1 \underline{a} f g), \Phi_2 \underline{a} f g, g(\Phi_1 \underline{a} f g, \Phi_2 \underline{a} f g), \underline{a})$$

holds in a suitable extension of **PL**.

We say (following Kreisel [15]) that Φ_1, Φ_2 satisfy the **no-counterexample interpretation** of A (short: Φ_1, Φ_2 *n.c.i.* A).

If A is no longer logically true but provable in some first-order theory, e.g. **PA**, then definition by cases will not be sufficient in general. In the case of **PA** for instance one needs all α -recursive functionals for $\alpha < \varepsilon_0$ and these functionals are also sufficient. This was proved firstly in [16] using an ε -substitution procedure based on [1].² Later Schwichtenberg [25] gave a proof of this result using a form of cut-elimination (due to [30]) instead.

The cut-elimination procedure does not give a local interpretation of proofs, i.e. given proofs of A and $A \rightarrow B$, a realization of the n.c.i. of B is not computed out of given realizations for the n.c.i of A and $A \rightarrow B$ but by a **global proof transformation** of the proof of B (which in general will cause a non-elementary increase in the length of this proof).³

The method of ε -substitution can be used (as indicated in the proof of the condition (δ) , to be discussed below, in [16]) to obtain $\beta (< \varepsilon_0)$ -recursive functionals satisfying the n.c.i. of B out of **given** $\alpha (< \varepsilon_0)$ -**recursive functionals** satisfying the n.c.i. of (prenex normal forms of) A and $A \rightarrow B$. This method however (which again in general has a non-elementary complexity in the logical depth of A) does not yield a uniform procedure (given by functionals of type level 3) which would provide functionals satisfying the n.c.i. of B **uniformly** in arbitrary functionals satisfying the n.c.i. of A and $A \rightarrow B$.

A third way to prove the no-counterexample interpretation of **PA** (by functionals which are $\alpha (< \varepsilon_0)$ -recursive) is via Gödel's functional interpretation (combined with negative translation) of **PA** in the calculus T of primitive recursive functionals of finite type (see e.g. [31]). This (combination of negative translation and) functional interpretation is a **local**

²A formalization of the method of ε -substitution was given by [29] and used in [20](thm.12).

³One should also mention here Gödel's discussion of Gentzen's 1936 consistency proof in his amazing 'Vortrag bei Zilsel' from 1938, first published (together with an English translation in [8]). Here Gödel interprets Gentzen's proof in terms of the no-counterexample interpretation and gives a discussion of the modus ponens rule in these terms which emphasizes the fact that this rule is decisive for the ordinal exponentiation indicating even a kind of local treatment of this rule, however without giving any details ([8] pp. 108-110). See also the illuminating remarks in [27].

proof transformation but at the level of the functional interpretation (of the negative translation) of A and $A \rightarrow B$ and not at the level of their n.c.i.: realizing functionals for $(B')^D$ can be obtained uniformly in any realizations of $(A')^D$ and $((A \rightarrow B)')^D$ by a simple typed lambda term (depending only on the logical form of A and B (Here A' and A^D denote the negative translation and the functional interpretation of A)).

The passage through higher types makes it necessary to use a normalization procedure for T in order to obtain the n.c.i. in terms of $\alpha(< \varepsilon_0)$ -recursive functionals rather than type 2 functionals defined in terms of primitive recursion in higher types (see e.g [21],[25]).

Instead of functional interpretation one could also use a combination of (negative translation plus) the Friedman-Dragalin A-translation and a suitable notion of realizability. If one uses here the so-called ‘minimal realizability’ of [3] one can avoid the use of higher types but the resulting interpretation again is not local at the level of the n.c.i. but only at the level of the ‘minimal realizability’ interpretation of (the Friedman-Dragalin translation of the negative translation of) $A, A \rightarrow B$.⁴

In this paper we calibrate the complexity of performing the modus ponens rule directly on the level of the n.c.i. without using higher types. It turns out that even for **PA**-provable sentences A and $A \rightarrow B$ with n.c.i. in T_0 no fixed subsystem T_n of T suffices:⁵ for every $n \in \mathbb{N}$ there are **PA**-provable sentence A, B ($B \in \Pi_2^0$) and functionals in T_0 satisfying (provably in $\widehat{\mathbf{PA}}^\omega \upharpoonright^6$) the n.c.i. of (arbitrary prenex normal forms of) A and $A \rightarrow B$ such that the n.c.i. of B is not satisfied by any function(al) $\in T_n$ (since with A and $A \rightarrow B$ also B is provable in **PA**, it is clear that the n.c.i. of B can be carried out in T). So already for **PA**-provable sentences the modus-ponens-complexity of the n.c.i. is not lower than the complexity of the n.c.i. of the whole theory **PA**. If A and $A \rightarrow B$ are not assumed to be provable in **PA**, then even T is not sufficient to solve the n.c.i. of the modus ponens rule (uniformly in functionals satisfying the n.c.i. of the assumptions) but $\widehat{\mathbf{PR}}^\omega + BR_{0,1}$ is, where $BR_{0,1}$ is the schema of bar recursion for bar recursion of type 0 (with values of type 1) and $\widehat{\mathbf{PR}}^\omega$ are all predicative primitive recursive functionals of finite type (in the sense of [13],[4]).

In special cases we can even solve the n.c.i. of the modus ponens as a **unification problem** yielding functionals satisfying the n.c.i. of B by unification (not depending on the quantifier-free part of A, B but only on the quantifier-prefix of their prenex normal forms): This is

⁴In connection with [3] one should mention that some of the result obtained in this paper by ‘minimal realizability’ can in fact be derived (sometimes in much stronger form) using only well-known facts from the literature ([31],[24]), see [14].

⁵ T_n denotes the fragment of Gödel’s T (see [7]) with R_ρ for $\text{deg}(\rho) \leq n$ only.

⁶ $\widehat{\mathbf{PA}}^\omega \upharpoonright$ is the subsystem of \mathbf{PA}^ω based on T_0 instead of T and with quantifier-free induction only, see [4] and section 2 below.

true for $A \in \Pi_3^0$ and $B \in \Pi_\infty^0$ (but already in this case T is not sufficient). This particular matter will be studied further in a subsequent paper.

Kreisel introduced his n.c.i. of arithmetic as an instance of his general definition of an ‘interpretation of a system Σ ’ which we recall here from [15]:

‘A computable function $f(n, a)$ is called an interpretation of the system Σ if

- (α) $f(n, a)$ is the number of a free variable formula A_n when a is the number of a formula \mathcal{A} of Σ (some Gödel numbering being assumed),
- (β) if \mathcal{A} is proved in Σ , from the proof we find an \bar{n} such that $A_{\bar{n}}$ is verifiable,
- (γ) if $\neg\mathcal{A}$ is proved in Σ , for each n we find a substitution for the (individual and function) variables of A_n which makes A_n false,
- (δ) if \mathcal{B} is proved from \mathcal{A} in Σ , we find a $g(n)$ so that $B_{g(n)}$ is verifiable if A_n is verifiable.’

For the n.c.i. of **PA** by $\alpha(< \varepsilon_0)$ -recursive functionals (resp. functionals in T) condition (α) follows immediately from the fact that the resulting set of free variable formulas is recursively enumerable. Condition (β) follows from each of the proof-methods discussed above. The condition (γ) and in particular the condition (δ) however (which are proved in [16] using the method of ε -substitution) do not follow from the approaches to the n.c.i. by cut-elimination or functional interpretation (or the Friedman-Dragalin translation plus realizability). The condition (δ) can be formulated in the case of the no-counterexample interpretation of **PA** in T (or, slightly reformulated, for $\alpha(< \varepsilon_0)$ -recursive functionals) as follows

$$(\delta) : \left\{ \begin{array}{l} \text{If } \underline{\Phi}_A \text{ n.c.i. } A \text{ is true for } \underline{\Phi}_A \in T \text{ and } \mathbf{PA} \vdash A \rightarrow B. \\ \text{Then one can construct } \underline{\Phi}_B \in T \text{ such that } \underline{\Phi}_B \text{ n.c.i. } B \text{ is true.} \end{array} \right.$$

Using (a careful analysis of the computational strength of) bar recursion of type 0 we give a new prove of Kreisel’s results including a strengthened uniform version of his condition (δ).

The condition (γ) translates in the case of the n.c.i. of **PA** at hand into

$$(\gamma) \left\{ \begin{array}{l} \text{If } A \equiv \exists x_1 \forall y_1 \dots \exists x_k \forall y_k A_0(x_1, y_1, \dots, x_k, y_k, \underline{a}) \in \mathcal{L}(\mathbf{PA}) \text{ and } \mathbf{PA} \vdash \neg A. \\ \text{Then constructively it holds that for all closed terms } \underline{\Phi} \in T \text{ (of suitable types)} \\ \text{there are } \underline{h} \text{ such that } A_0(\Phi_1 \underline{h}, h_1(\Phi_1 \underline{h}), \dots, \Phi_k \underline{h}, h_k(\Phi_1 \underline{h}, \dots, \Phi_k \underline{h}), \underline{a}) \text{ is false.} \end{array} \right.$$

Classically the existence of \underline{h} satisfying (γ) can be shown quite easily (see remark 4.10). A constructive proof of (γ) was given in [16], again by the use of the ε -substitution method. We give a new proof of a uniform strengthening of (γ) in section 4.

Finally we discuss a different definition of interpretation presented in [17] and show that this definition is incomparable with the definition given in [15]. In particular we show: the n.c.i. of \mathbf{PA}_{n+1} (the fragment of \mathbf{PA} with Σ_{n+1}^0 -induction only) in T_n (which holds by [22]) is an interpretation in the sense of [17] but not in the sense of [15] since the condition (δ) is violated in this case.

2 The modus ponens complexity of the no-counterexample interpretation for PA-provable sentences

Definition 2.1 Let $A := \exists x_1 \forall y_1 \dots \exists x_k \forall y_k A_0(x_1, y_1, \dots, x_k, y_k, \underline{a})^7$ be a formula in the language $\mathcal{L}(\mathbf{PA})$ of Peano arithmetic \mathbf{PA} (which for convenience is assumed to contain symbols for every primitive recursive function with the corresponding defining equations as axiom of \mathbf{PA}).

The Herbrand normal form A^H of A is defined by

$$A^H := \forall h_1, \dots, h_k \exists x_1, \dots, x_k \overbrace{A_0(x_1, h_1 x_1, \dots, x_k, h_k x_1 \dots x_k, \underline{a})}^{A_0^H :=}$$

A tuple $\underline{\Phi} (= \Phi_1, \dots, \Phi_k)$ of functionals of type levels ≤ 2 satisfies the no-counterexample interpretation of A if $\underline{\Phi} \underline{a} \underline{h}$ realizes $\exists \underline{x}$ ' (where $\underline{h} := h_1, \dots, h_k$ and $\underline{x} := x_1, \dots, x_k$), i.e. if

$$\forall \underline{a}, \underline{h} A_0(\Phi_1 \underline{a} \underline{h}, h_1(\Phi_1 \underline{a} \underline{h}), \dots, \Phi_k \underline{a} \underline{h}, h_k(\Phi_1 \underline{a} \underline{h}, \dots, \Phi_k \underline{a} \underline{h}), \underline{a}).$$

In this case we write ' $\underline{\Phi}$ n.c.i. A '.

In the following \mathbf{PRA} denotes primitive recursive arithmetic extended by classical first-order predicate logic. \mathbf{PA}^ω (resp. \mathbf{HA}^ω) is the classical (resp. intuitionistic) arithmetic in all finite types with full induction and all primitive recursive functionals in the sense of Gödel and a quantifier-free rule of extensionality (so in the terminology of [31], \mathbf{HA}^ω is the system $\mathbf{WE-HA}^\omega$). $\widehat{\mathbf{PA}}^\omega \upharpoonright$ (resp. $\widehat{\mathbf{HA}}^\omega \upharpoonright$) denotes the fragment of \mathbf{PA}^ω (resp. \mathbf{HA}^ω) with quantifier-free induction only and the Gödel-recursors R_ρ replaced by the predicative Kleene-recursors \widehat{R}_ρ (this systems was introduced and studied in [4]). By T and $\widehat{\mathbf{PR}}^\omega$ we denote the quantifier-free parts (in the sense on [31](1.6.13)) of \mathbf{PA}^ω and $\widehat{\mathbf{PA}}^\omega \upharpoonright$ respectively.

⁷Here \underline{a} are all the free variables of A .

T_n is the fragment of T with R_ρ for ρ of level $\leq n$ only. $\widehat{\mathbf{PR}}^\omega$ is simply a definitorial extension of T_0 since $R_0 = \widehat{R}_0$ and \widehat{R}_ρ for $\rho > 0$ is definable from \widehat{R}_0 by λ -abstraction.

The type level or degree $\text{deg}(\rho)$ of a type ρ is defined as $\text{deg}(0) := 0$, $\text{deg}(\rho(\tau)) := \max(\text{deg}(\tau) + 1, \text{deg}(\rho))$.

Convention: By the phrase ‘a functional $\Phi \in T_{(n)}$ ’ we always mean ‘a closed term Φ of $T_{(n)}$ ’. Sometimes we only write $\Phi \in T_{(n)}$ but again always refer to a closed term of $T_{(n)}$ representing the functional.

Proposition 2.2 *For every $n \in \mathbb{N}$ there are sentence (i.e. closed formulas) A, B such that*

- 1) A is prenex,
- 2) $B \equiv \forall x \exists y B_0(x, y) \in \Pi_2^0$,
- 3) $\mathbf{PRA} \vdash A$,
- 4) $\mathbf{PA} \vdash A \rightarrow B$,
- 5) A as well as every prenex normal form $(A \rightarrow B)^{pr}$ of $A \rightarrow B$ has (provably in $\widehat{\mathbf{HA}}^\omega \upharpoonright$) a n.c.i. by suitable functionals in T_0 , i.e.

$$\widehat{\mathbf{HA}}^\omega \upharpoonright \vdash \underline{\Phi}_A \text{ n.c.i. } A \wedge \underline{\Phi}_{(A \rightarrow B)^{pr}} \text{ n.c.i. } (A \rightarrow B)^{pr}$$

with $\underline{\Phi}_A, \underline{\Phi}_{(A \rightarrow B)^{pr}} \in T_0$,

but:

- 6) there is no function $\varphi \in T_n$ which satisfies the n.c.i. of B , i.e. there is no $\varphi \in T_n$ for which $\forall x B_0(x, \varphi x)$ is true in the standard model of \mathbf{PA} .

Proof: Let $n \in \mathbb{N}$ be fixed. It is well-known that the provably recursive functions of \mathbf{PA} are just the $\alpha(< \varepsilon_0)$ -recursive functions. Since the definable functions of type 1 in T_n are $< \omega_{n+1}(\omega)$ -recursive (see [21]), there is a Π_2^0 -sentence $B \equiv \forall x \exists y B_0(x, y)$ in $\mathcal{L}(\mathbf{PA})$ (namely $\forall x \exists y T(\bar{e}, x, y)$ for a certain numeral \bar{e}) such that $\mathbf{PA} \vdash B$, but there is no $t^1 \in T_n$ for which $\forall x \in \mathbb{N} B_0(x, tx)$ is true.

Since $\mathbf{PA} \vdash B$ there are finitely many instances $\tilde{A}_1, \dots, \tilde{A}_k$ such that for their universal closures A_1, \dots, A_k

$$\mathbf{PRA} \vdash \bigwedge_{i=1}^k A_i \rightarrow B.$$

Let $\widehat{A}_i(x, \underline{a})$ be the induction formula corresponding to A_i , where x is the induction variable and \underline{a} includes all parameters, i.e.

$$A_i \leftrightarrow \forall \underline{a} (\widehat{A}_i(\mathbf{0}, \underline{a}) \wedge \forall x (\widehat{A}_i(x, \underline{a}) \rightarrow \widehat{A}_i(x', \underline{a})) \rightarrow \forall x \widehat{A}_i(x, \underline{a})).$$

We now define

$$A := \forall x, \underline{a} \exists y_1, \dots, y_k \bigwedge_{i=1}^k (y_i = \mathbf{0} \leftrightarrow \widehat{A}_i(x, \underline{a})).$$

It is clear that

- (i) **PRA** $\vdash A$ (in fact predicate logic with equality plus the axiom $0 \neq S0$ suffices),
- (ii) **PA** $\vdash A \rightarrow B$.

In **PRA**, the variables $x_1, \dots, x_k, \underline{a}$ and the variables y_1, \dots, y_k can be coded together as single variables x, y . Although we do not carry out this coding for the sake of better readability we are free to consider these tuples as single variables from now on. As a consequence we only have to deal with the following prenex normal forms of $A \rightarrow \forall u \exists v B_0(u, v)$

$$(1) \exists x, \underline{a} \forall u \exists v \forall \underline{y} (\dots)^{pr},$$

$$(2) \exists x, \underline{a} \forall \underline{y} (\dots)^{pr},$$

$$(3) \forall u \exists x, \underline{a} \forall \underline{y} (\dots)^{pr},$$

$$(4) \forall u \exists v, x, \underline{a} \forall \underline{y} (\dots)^{pr},$$

where $(\dots)^{pr}$ refer to any prenex normal form of the remaining formula in each case.

For $i = 1, \dots, 4$ the Herbrand normal form $(i)^H$ of (i) is implied by the partial Herbrand normal form where Herbrand index functions are introduced only for the universal quantifiers in front of $(\dots)^{pr}$. So e.g. for (1), $(1)^H$ is implied by

$$(+) \forall f, \underline{g} \exists x, \underline{a}, v ([u/f(x, \underline{a})], [\underline{y}/\underline{g}(x, \underline{a}, v)])^{pr}.$$

One easily shows by classical logic (and λ -abstraction) that $(+)$ is equivalent to

$$(*) \exists \underline{g} \forall x, \underline{a} \left(\bigwedge_{i=1}^k (g_i x \underline{a} = \mathbf{0} \leftrightarrow A_i(x, \underline{a})) \rightarrow \forall u \exists v B_0(u, v) \right).$$

In fact

$$\begin{aligned}
(+)&\Leftrightarrow \\
&\forall f, \underline{g}(\forall x, \underline{a}, v \bigwedge_{i=1}^k (g_i x \underline{a} v = 0 \leftrightarrow \widehat{A}_i(x, \underline{a})) \rightarrow \exists x, \underline{a}, v B_0(f x \underline{a}, v)) \Leftrightarrow \\
&\exists \underline{g} \forall x, \underline{a}, v \bigwedge_{i=1}^k (g_i x \underline{a} v = 0 \leftrightarrow \widehat{A}_i(x, \underline{a})) \rightarrow \forall f \exists x, \underline{a}, v B_0(f x \underline{a}, v) \Leftrightarrow \\
&\exists \underline{g} \forall x, \underline{a}, v \bigwedge_{i=1}^k (g_i x \underline{a} v = 0 \leftrightarrow \widehat{A}_i(x, \underline{a})) \rightarrow \forall u \exists v B_0(u, v) \Leftrightarrow \\
&\exists \underline{g} \forall x, \underline{a} \bigwedge_{i=1}^k (g_i x \underline{a} = 0 \leftrightarrow \widehat{A}_i(x, \underline{a})) \rightarrow \forall u \exists v B_0(u, v).
\end{aligned}$$

In a similar way one shows the corresponding result for (2), (3), (4). So put together we have

$$(*) \rightarrow (i)^H, \text{ where } i = 1, \dots, 4,$$

by predicate logic (and λ -abstraction). But $(*)$ and therefore $(i)^H$ is provable in \mathbf{PRA}^2 which is the extension of \mathbf{PRA} by adding function quantifiers to \mathbf{PRA} and allowing function variables to occur in the schema of quantifier-free induction

$$\mathbf{QF-IA} : A_0(0) \wedge \forall x(A_0(x) \rightarrow A_0(x')) \rightarrow \forall x A_0(x).$$

This follows simply by applying $\mathbf{QF-IA}$ to $A_0(x) := (g_i x \underline{a} = 0)$ which yields A_i and so $\forall u \exists v B_0(u, v)$.

So $\mathbf{PRA}^2 \vdash ((A \rightarrow B)^{pr})^H$ for every prenex normal form of $A \rightarrow B$.

However \mathbf{PRA}^2 has (via negative translation) a functional interpretation and hence a n.c.i. in $\widehat{\mathbf{HA}}^\omega \upharpoonright$ by terms $\in T_0$. Thus there are functionals $\underline{\Phi}_{(A \rightarrow B)^{pr}} \in T_0$ such that

$$\widehat{\mathbf{HA}}^\omega \upharpoonright \vdash \underline{\Phi}_{(A \rightarrow B)^{pr}} \text{ n.c.i. } (A \rightarrow B)^{pr}$$

for each prenex normal form of $A \rightarrow B$. The same holds true for A which is even provable in \mathbf{PRA} : there are functionals $\underline{\Phi}_A \in T_0$ such that

$$\widehat{\mathbf{HA}}^\omega \upharpoonright \vdash \underline{\Phi}_A \text{ n.c.i. } A,$$

which concludes the proof of the proposition. \square

Remark 2.3 We can replace ‘ $\varphi \in T_n$ ’, ‘ $\underline{\Phi} \in T_0$ ’ in the proposition above by ‘ φ is $\alpha(< \omega_{n+1}(\omega))$ -recursive’ and ‘ $\underline{\Phi}$ is primitive recursive in the sense of Kleene’, since the closed terms $t^2 \in T_n$ denote just the $\alpha(< \omega_{n+1}(\omega))$ -recursive functionals (see e.g. [21]). In the following we only state the T_n -versions of our results explicitly since it is straightforward to formulate them in terms of ordinal recursive function(al)s as well.

We now consider the condition (δ) mentioned in the introduction. This condition was verified for the n.c.i. of \mathbf{PA} (by $\alpha(< \varepsilon_0)$ -recursive functionals) in [16] using the method of ε -substitution. It does not follow from the proofs of the n.c.i. by cut-elimination or functional interpretation. In section 4 below we will prove a new strong uniform version of this condition.

Let \mathbf{PA}_n be the subsystem of \mathbf{PA} with induction restricted to Σ_n^0 -formulas. In [22] it is shown that \mathbf{PA}_{n+1} has (via negative translation) a functional interpretation in T_n . Hence also the n.c.i. of \mathbf{PA}_{n+1} -provable formulas can be satisfied in T_n . However as a corollary of proposition 2.2 we have

Corollary 2.4 *The no-counterexample interpretation of \mathbf{PA}_{n+1} in T_n (or – equivalently – by $\alpha(< \omega_{n+1}(\omega))$ -recursive functionals) does not satisfy the condition (δ) and hence is not an interpretation in the sense of [15].*

Proof: Choose $A, B \in \mathcal{L}(\mathbf{PA}_{n+1})$ as in proposition 2.2 and let $(A \rightarrow B)^{pr}$ be any prenex normal form of $A \rightarrow B$ and \tilde{A} be the prenex normal form of $A \wedge (A \rightarrow B)^{pr}$ which results e.g. by shifting first all A -quantifiers to the front and then all $(A \rightarrow B)^{pr}$ -quantifiers. Already by classical logic, \tilde{A} implies B and so in particular

$$\mathbf{PA}_{n+1} \vdash \tilde{A} \rightarrow B.$$

From proposition 2.2 it follows that both A and $(A \rightarrow B)^{pr}$ have a n.c.i. by functionals in T_0 (i.e. by $\alpha(< \omega^\omega)$ -recursive and hence ordinary primitive recursive functionals). From this one easily constructs functionals in T_0 satisfying the n.c.i. of \tilde{A} . However, again by proposition 2.2, B does not have a n.c.i. in T_n (and hence not by an $\alpha(< \omega_{n+1}(\omega))$ -recursive function). So \tilde{A} and B provide a counterexample to the condition (δ) for the n.c.i. of \mathbf{PA}_{n+1} in T_n . \square

3 The uniform modus ponens complexity of the no-counterexample interpretation for arbitrary formulas $A, B \in \mathcal{L}(\mathbf{PA})$

Definition 3.1 *A pair $(\mathcal{T}, \mathcal{F})$ consisting of a theory \mathcal{T} and a quantifier-free functional calculus $\mathcal{F} \subset \mathcal{T}$ suffices for the uniform n.c.i. of the modus ponens rule if for all (prenex) formulas $A, B \in \Pi_\infty^0$ ($A, B \in \mathcal{L}(\mathbf{PA})$) and every prenex normal form $(A \rightarrow B)^{pr}$ of $A \rightarrow B$ there are functionals $\underline{\Psi} \in \mathcal{F}$ (i.e. closed terms $\underline{\Psi}$ of \mathcal{F}) such that*

$$\begin{aligned} \mathcal{T} \vdash \forall \underline{\Phi}_A, \underline{\Phi}_{(A \rightarrow B)^{pr}} \left((\underline{\Phi}_A \text{ n.c.i. } A) \wedge (\underline{\Phi}_{(A \rightarrow B)^{pr}} \text{ n.c.i. } (A \rightarrow B)^{pr}) \right. \\ \left. \rightarrow \underline{\Psi}(\underline{\Phi}_A, \underline{\Phi}_{(A \rightarrow B)^{pr}}) \text{ n.c.i. } B \right). \end{aligned}$$

Proposition 3.2 *There are sentences $A, B \in \mathcal{L}(\mathbf{PA})$ in prenex normal form such that for all prenex normal forms $(A \rightarrow B)^{pr}$ of $A \rightarrow B$*

$$\mathbf{PA}^\omega \not\vdash \exists \underline{\Phi}_A (\underline{\Phi}_A \text{ n.c.i. } A) \wedge \exists \underline{\Phi}_{(A \rightarrow B)^{pr}} (\underline{\Phi}_{(A \rightarrow B)^{pr}} \text{ n.c.i. } (A \rightarrow B)^{pr}) \rightarrow \exists \underline{\Phi}_B (\underline{\Phi}_B \text{ n.c.i. } B).$$

Moreover we can take $A \in \Pi_3^0$ and B quantifier-free (so that $(\underline{\Phi}_B \text{ n.c.i. } B) \leftrightarrow B$ with $\underline{\Phi}_B$ being the empty tuple).

Proof: Let $A := \forall x \exists y \forall z (Txy \vee \neg Txxz)$, where T denotes Kleene's T -predicate, and $B := (0 = 1)$. There is only one prenex normal form of $A \rightarrow B$:

$$\exists x \forall y \exists z (Txy \vee \neg Txxz \rightarrow 0 = 1)$$

and its n.c.i. requires a functionals Φ_1, Φ_2 such that

$$(*) \forall f (T(\Phi_1 f, \Phi_1 f, f(\Phi_1 f)) \vee \neg T(\Phi_1 f, \Phi_1 f, \Phi_2 f) \rightarrow 0 = 1).$$

The n.c.i. of A is realized by a functional Φ_0 such that

$$(**) \forall x, g (T(x, x, \Phi_0 xg) \vee \neg T(x, x, g(\Phi_0 xg))).$$

We now show that

$$\mathbf{PA}^\omega \not\vdash \exists \Phi_0, \Phi_1, \Phi_2 ((*) \wedge (**)) \rightarrow 0 = 1.$$

We have to show that $\mathbf{PA}^\omega + \exists \Phi_0, \Phi_1, \Phi_2 ((*) \wedge (**))$ is consistent:

$$\text{Define } \Phi_0 xg := \begin{cases} g0, & \text{if } T(x, x, g0) \\ 0, & \text{otherwise.} \end{cases}$$

Then one easily verifies that

$$\mathbf{PA}^\omega \vdash \forall x, g (T(x, x, \Phi_0 xg) \vee \neg T(x, x, g(\Phi_0 xg))).$$

Next we show that

$$\begin{aligned} & \mathbf{PA}^\omega + \text{AC}^{1,0}\text{-qf} + \forall f^1 (f \text{ is recursive}) \vdash \\ & \exists \Phi_1, \Phi_2 \forall f (T(\Phi_1 f, \Phi_1 f, f(\Phi_1 f)) \vee \neg T(\Phi_1 f, \Phi_1 f, \Phi_2 f) \rightarrow 0 = 1). \end{aligned}$$

This however follows from the fact that $\mathbf{PA}^\omega + \forall f^1 (f \text{ is recursive})$ proves (using the undecidability of the halting problem)

$$\forall f \exists x, z (T(x, x, fx) \vee \neg T(x, x, z) \rightarrow 0 = 1),$$

which implies using $\text{AC}^{1,0}\text{-qf}$

$$\exists \Phi_1, \Phi_2 \forall f (T(\Phi_1 f, \Phi_1 f, f(\Phi_1 f)) \vee \neg T(\Phi_1 f, \Phi_1 f, \Phi_2 f) \rightarrow 0 = 1).$$

The proof is now finished by verifying the consistency of $\mathbf{PA}^\omega + \text{AC}^{1,0}\text{-qf} + \forall f^1 (f \text{ is recursive})$ which however follows from the fact that

$$\mathbf{HEO} \models \mathbf{PA}^\omega + \text{AC}^{1,0}\text{-qf} + \forall f^1 (f \text{ is recursive}),$$

where \mathbf{HEO} is the type structure of the hereditarily effective operations in all finite types (the fact that \mathbf{HEO} forms a model of \mathbf{PA}^ω is proved in [31]. That it is a model of $\text{AC}^{1,0}\text{-qf}$ follows from the fact that one can always find an effective choice functional by unbounded search since quantifier-free formulas of \mathbf{PA}^ω are decidable). \square

Corollary 3.3 (\mathbf{PA}^ω, T) does not suffice for the uniform n.c.i. of the modus ponens rule.

Remark 3.4 1) The proof above does not exclude the possibility that e.g.

$(\mathbf{PA}^\omega + \text{AC}_{ar}^{0,0}, T)$ satisfies the uniform n.c.i. of the modus ponens rule which remains an open problem. Nevertheless we will show below that even (\mathcal{S}^ω, T) does not suffice to solve uniformly the unification problem associated with the n.c.i. of the modus ponens (which however does not exclude other ways of satisfying the n.c.i. of the modus ponens)

2) In section 4 below we will show that (\mathbf{PA}^ω, T) suffices for the **pointwise** n.c.i. of the modus ponens in the sense that one can construct functionals of type level 3 in a **genuine extension** of T which produce out of given functionals $\in T$ which satisfy the n.c.i. of $A, (A \rightarrow B)^{pr}$ functionals $\in T$ which satisfy the n.c.i. of B .

We now show that both $(\widehat{\mathbf{HA}}^\omega \upharpoonright + \mu, \widehat{\mathbf{PR}}^\omega + \mu)$ and $(\widehat{\mathbf{HA}}^\omega \upharpoonright + \text{BR}_{0,1}, \widehat{\mathbf{PR}}^\omega + \text{BR}_{0,1})$ do suffice:

Definition 3.5 ([4]) $\mathcal{T}^\omega + \mu$ is the extension of \mathcal{T}^ω obtained by adding a constant μ^2 together with the axioms

$$(\mu) : f^1 x =_0 0 \rightarrow f(\mu f) =_0 0 \wedge \mu f \leq_0 x, f(\mu f) \neq_0 \rightarrow \mu f =_0 0.$$

Definition 3.6 ([28]) $\mathcal{T}^\omega + \text{BR}_{\rho,\tau}$ is the extension of \mathcal{T}^ω obtained by adding the bar recursion constant $B_{\rho,\tau}$ with the axioms

$$(\text{BR}_{\rho,\tau}) : \begin{cases} x(\overline{y, n^0}) < n \rightarrow B_{\rho,\tau} x z u n y =_\tau z n y \\ x(\overline{y, n}) \geq n \rightarrow B_{\rho,\tau} x z u n y =_\tau u(\lambda D^\rho. B_{\rho,\tau} x z u n'(\overline{y, n} * D)) n y, \end{cases}$$

where y is of type $0(\rho 0)$ and u is of type $\tau(\rho 0)(0)(\tau\rho)$ and

$$(\overline{y}, \overline{n} * D)(k^0) =_\rho \begin{cases} yk, & \text{if } k < n \\ D, & \text{if } k = n \\ 0^\rho, & \text{otherwise.} \end{cases}$$

Proposition 3.7 *Let $(\mathcal{T}, \mathcal{F})$ be either $(\widehat{\mathbf{HA}}^\omega \upharpoonright + \mu, \widehat{\mathbf{PR}}^\omega + \mu)$ or $(\widehat{\mathbf{HA}}^\omega \upharpoonright + BR_{0,1}, \widehat{\mathbf{PR}}^\omega + BR_{0,1})$. Then $(\mathcal{T}, \mathcal{F})$ suffices for the n.c.i. of the modus ponens (uniformly in functionals satisfying the n.c.i. any of prenex normal forms of A and $A \rightarrow B$).*

Proof: Lets consider the schema of arithmetical choice

$$\text{AC}_{ar}^{0,0} : \forall x \exists y A(x, y) \rightarrow \exists f \forall x, y A(x, fx),$$

where $A \in \Pi_\infty^0$ (A may contain function parameters).

One easily verifies that

$$\widehat{\mathbf{PA}}^\omega \upharpoonright + \text{AC}_{ar}^{0,0} \vdash (\exists \Phi_A (\Phi_A \text{ n.c.i. } A)) \rightarrow A$$

for all prenex formulas $A \in \Pi_\infty^0$. Since furthermore

$$\widehat{\mathbf{PA}}^\omega \upharpoonright \vdash B \rightarrow B^H$$

for all prenex formulas $B \in \Pi_\infty^0$ we have

$$\begin{aligned} & \widehat{\mathbf{PA}}^\omega \upharpoonright + \text{AC}_{ar}^{0,0} \vdash \\ & \exists \Phi_A (\Phi_A \text{ n.c.i. } A) \wedge \exists \Phi_{(A \rightarrow B)^{pr}} (\Phi_{(A \rightarrow B)^{pr}} \text{ n.c.i. } (A \rightarrow B)^{pr}) \rightarrow \forall \underline{h}, \underline{a} \exists \underline{x} B_0^H(\underline{h}, \underline{x}, \underline{a}), \end{aligned}$$

where $\forall \underline{h} \exists \underline{x} B_0^H(\underline{h}, \underline{x}, \underline{a})$ is the Herbrand normal form $B^H(\underline{a})$ of $B(\underline{a})$ and \underline{a} are all free variables of B .

$\widehat{\mathbf{PA}}^\omega \upharpoonright + \text{AC}_{ar}^{0,0}$ has (via negative translation) a functional interpretation in \mathcal{T} by terms $\in \mathcal{F}$. For $(\widehat{\mathbf{HA}}^\omega \upharpoonright + \mu, \widehat{\mathbf{PR}}^\omega + \mu)$ this is proved in [4]. For $(\widehat{\mathbf{HA}}^\omega \upharpoonright + BR_{0,1}, \widehat{\mathbf{PR}}^\omega + BR_{0,1})$ this follows from [28] using the facts that $\widehat{\mathbf{PA}}^\omega \upharpoonright$ has an interpretation in $\widehat{\mathbf{PR}}^\omega$, that $\text{AC}_{ar}^{0,0}$ is derivable in $\widehat{\mathbf{PA}}^\omega \upharpoonright + \Pi_1^0\text{-AC}^{0,0}$ (note that $\widehat{\mathbf{PA}}^\omega \upharpoonright + \Pi_1^0\text{-AC}^{0,0} \vdash \Pi_1^0\text{-CA}$ and so by iteration – using the presence of function parameters in $\Pi_1^0\text{-CA}$ – also $\widehat{\mathbf{PA}}^\omega \upharpoonright + \Pi_1^0\text{-AC}^{0,0} \vdash \Pi_\infty^0\text{-CA}$ and therefore $\widehat{\mathbf{PA}}^\omega \upharpoonright + \Pi_1^0\text{-AC}^{0,0} \vdash \text{AC}_{ar}^{0,0}$) and that the interpretation of $\Pi_1^0\text{-AC}^{0,0}$ uses only $B_{0,1}$ and functionals from $\widehat{\mathbf{PR}}^\omega$. Note that the crucial lemma 1 from [28] (restricted to $B_{0,1}$) can easily be proved in $\widehat{\mathbf{PR}}^\omega + BR_{0,1}$.

Hence there are functionals $\tilde{\Psi} \in \mathcal{F}$ such that

$$(+)\left\{\begin{array}{l} \mathcal{T} \vdash \forall \underline{\Phi}_A, \underline{\Phi}_{(A \rightarrow B)^{pr}} \left((\underline{\Phi}_A \text{ n.c.i. } A) \wedge (\underline{\Phi}_{(A \rightarrow B)^{pr}} \text{ n.c.i. } (A \rightarrow B)^{pr}) \right. \\ \left. \rightarrow \forall \underline{h}, \underline{a} B_0^H(\underline{h}, \tilde{\Psi}(\underline{\Phi}_A, \underline{\Phi}_{(A \rightarrow B)^{pr}}, \underline{h}, \underline{a}), \underline{a}) \right). \end{array}\right.$$

Thus $\underline{\Psi} := \lambda \underline{h}. \tilde{\Psi}(\underline{\Phi}_A, \underline{\Phi}_{(A \rightarrow B)^{pr}}, \underline{h}, \underline{a})$ satisfies the claim made in the proposition. \square

Remark 3.8 1) Similar to $\underline{\Psi}$ one can also extract $\underline{\xi}, \underline{\zeta} \in \widehat{\mathbf{PR}}^\omega + BR_{0,1}$ realizing the universal function quantifiers hidden in ' $\underline{\Phi}_A$ n.c.i. A ' and ' $\underline{\Phi}_{(A \rightarrow B)^{pr}}$ n.c.i. $(A \rightarrow B)^{pr}$ '.

2) In the above proof, (+) can actually be strengthened by not assuming that $\underline{\Phi}_A$ (resp. $\underline{\Phi}_{(A \rightarrow B)^{pr}}$) satisfies the no-counterexample interpretation uniformly in the parameters \underline{a} of A^8 , i.e. we can quantify \underline{a} outside the whole implication in (+) and weaken ($\underline{\Phi}_A$ n.c.i. A) (and likewise also ($\underline{\Phi}_{(A \rightarrow B)^{pr}}$ n.c.i. $(A \rightarrow B)^{pr}$)) to

$$\forall \underline{h} A_0(\underline{\Phi}_1^A \underline{h}, h_1(\underline{\Phi}_1^A \underline{h}), \dots, \underline{\Phi}_k^A \underline{h}, h_k(\underline{\Phi}_1^A \underline{h}), \dots, \underline{\Phi}_k^A \underline{h}), \underline{a}.$$

I.e. we only require $\underline{\Phi}^A$ to satisfy the n.c.i. of A for the fixed parameters \underline{a} . As in the proof above we now obtain functionals $\underline{\chi}$ which satisfy the modus ponens uniformly in $\underline{h}, \underline{a}$ and functionals $\underline{\Phi}^A, \underline{\Phi}^{(A \rightarrow B)^{pr}}$ satisfying the n.c.i. for the parameters \underline{a} .

Corollary to the proof of proposition 3.7: The proof above immediately generalizes to the case where A and B contain function parameters $\underline{\alpha}, \underline{\beta}$ and yields functionals $\underline{\Psi}(\underline{\Phi}_A, \underline{\Phi}_{(A \rightarrow B)^{pr}}, \underline{\alpha}, \underline{\beta})$ which solve the corresponding modus ponens instance uniformly in $\underline{\Phi}_A, \underline{\Phi}_{(A \rightarrow B)^{pr}}$ and $\underline{\alpha}, \underline{\beta}$. This in particular implies that we can solve the modus ponens problem uniformly in arbitrary formulas A, B in $\mathcal{L}(\mathbf{PA})$ of **fixed** quantifier complexities since all formulas $A \in \Pi_n^0$ can be obtained from $\forall x_1 \exists y_1 \dots \forall x_n \exists y_n (\alpha(\underline{x}, \underline{y}) =_0 0)$ by substituting the characteristic function of the quantifier-free matrix of A (which can be defined in $\widehat{\mathbf{PR}}^\omega$) for the function variable α .

For $A \in \Pi_3^0, B \in \Pi_\infty^0$, the functionals $\underline{\Psi} \in \widehat{\mathbf{PR}}^\omega + BR_{0,1}$ solving the n.c.i. of the modus ponens rule (which exist by 3.7) can be obtained as the solution of a system of functional equations:

⁸Lets assume here for simplicity that A and B contain the same parameters \underline{a} . This can be achieved by introducing dummy variables if necessary.

Let $A := \forall x \exists y \forall z A_0(x, y, z)$ and $B := \forall u_1 \exists v_1 B_0(u_1, v_1, \dots)$. Consider the following prenex normal form of $A \rightarrow B$

$$(A \rightarrow B)^{pr} := \exists x \forall y \exists z \forall u_1 \exists v_1 \dots (A_0(x, y, z) \rightarrow B_0(u_1, v_1, \dots)).$$

Then

$$((A \rightarrow B)^{pr})^H \equiv \forall f, \underline{h} \exists x, z, \underline{v} (A_0(x, f\underline{h}, z) \rightarrow B_0(h_0xz, v_1, h_1xzv_1, v_2, \dots)).$$

So the n.c.i. of $(A \rightarrow B)^{pr}$ requires functionals $\Phi_1, \Phi_2, \Psi_1, \Psi_2, \dots$ such that

$$(*) \forall f, \underline{h} (A_0(\Phi_1 f \underline{h}, f(\Phi_1 f \underline{h}), \Phi_2 f \underline{h}) \rightarrow B_0(h_0(\Phi_1 f \underline{h}, \Phi_2 f \underline{h}), \Psi_1 f \underline{h}, \dots)).$$

Since

$$A^H \equiv \forall g, x \exists y A_0(x, y, gy),$$

the n.c.i. of A requires a functional Φ_0 such that

$$(**) \forall g, x A_0(x, \Phi_0 gx, g(\Phi_0 gx)).$$

To perform a modus ponens using $(*)$, $(**)$ to obtain a solution for the for the n.c.i. of B we solve the following systems of equations (mp-unification) for x, f, g (uniformly in $\underline{h}, \Phi_0, \Phi_1, \Phi_2$):

$$(1) \begin{cases} x =_0 \Phi_1 f \underline{h} \\ f(\Phi_1 f \underline{h}) =_0 \Phi_0 gx \\ \Phi_2 f \underline{h} =_0 g(\Phi_0 gx). \end{cases}$$

Let $f[\underline{h}, \underline{\Phi}]$ be the f -solution for $\underline{h}, \Phi_0, \Phi_1, \Phi_2$. Taking then $\tilde{h}_0 := \lambda x, y, u, \tilde{h}_i xyv_1 \dots v_i := h_i v_1 \dots v_i$ (for $i \geq 1$) and $\tilde{\Psi}_i(u, h_1, \dots) := \Psi_i(f[\tilde{h}, \underline{\Phi}], \tilde{h})$ we obtain that $\tilde{\Psi}$ n.c.i. B .

Remark 3.9 Note that the system of equation (1) is the same as the one resulting from the functional interpretation of the double-negation shift

$$\forall x^0 \neg \neg \exists y^0 \forall z^0 A_0(x, y, z) \rightarrow \neg \neg \forall x \exists y \forall z A_0(x, y, z)$$

solved by Spector [28] using bar recursion in his functional interpretation of classical analysis (via negative translation). For completeness we include here the solution.

In our case it suffices in fact to construct an f such that there exists a g so that (1) holds for $x = \Phi_1 f \underline{h}$, since the functionals $\tilde{\Psi}$ do not depend on g .

In fact we solve (following Spector [28])

$$(2) \exists f \forall n \leq \Phi_1 f \underline{h} \exists g_n (\Phi_0(g_n, n) =_0 f n \wedge g_n(f n) =_0 \Phi_2 f \underline{h})$$

for f . Note that this solves (1) as well: take $x := n := \Phi_1 f \underline{h}$ and $g := g_x$.

Solution of (2): Define

$$A(f, n) := n \leq \Phi_1 f \underline{h} \rightarrow \exists g_n (\Phi_0(g_n, n) = f n \wedge g_n(f n) = \Phi_2 f \underline{h}).$$

We define a functional $B \in 1(0)(1)$ which satisfies

- (i) $\forall i < x (B(\overline{f, x}; x)(i) = fi)$,
- (ii) $\forall n \geq x A(B(\overline{f, x}; x), n)$.

Then $B(0^1, 0^0)$ satisfies $\forall n A(B(0, 0), n)$, i.e. solves ‘ $\exists f$ ’ in (2).

We now define $B(\overline{f, x}; x)$ by bar recursion:

Case 1): $\Phi_1(\overline{f, x}) \underline{h} < x$. Take $B(\overline{f, x}; x) := \overline{f, x}$. Then $B(\overline{f, x}; x)$ trivially satisfies (i) and because of $n \geq x \rightarrow n > \Phi_1(\overline{f, x}) \underline{h}$ (by the case) also (ii).

Case 2): $\Phi_1(\overline{f, x}) \underline{h} \geq x$. By assumption $B(\overline{f, x * \langle X \rangle}; x')$ is defined already such that

- (i)' $\forall i \leq x (B(\overline{f, x * \langle X \rangle}; x')(i) = (\overline{f, x * \langle X \rangle})(i))$ and
- (ii)' $\forall n \geq x' A(B(\overline{f, x * \langle X \rangle}; x'), n)$ for all X (Note that $\overline{f, x * \langle X \rangle} = \overline{\overline{f, x * \langle X \rangle}, x'}$).

Define $B(\overline{f, x}; x) := B(\overline{f, x * \langle K \rangle}; x')$,

where $K := \Phi_0 g_x x$ and $g_x := \lambda X. \Phi_2(B(\overline{f, x * \langle X \rangle}; x')) \underline{h}$.

By (i)', (ii)' we have

$$\forall n \geq x' A(B(\overline{f, x}; x), n) \text{ and } \forall i < x (B(\overline{f, x}; x)(i) = fi).$$

So it remains to show $A(B(\overline{f, x}; x), x)$, i.e.

$$\exists g_x (\Phi_0(g_x, x) = B(\overline{f, x}; x)(x) \wedge g_x(B(\overline{f, x}; x)(x)) = \Phi_2(B(\overline{f, x}; x)) \underline{h}) :$$

$$B(\overline{f, x}; x)(x) = B(\overline{f, x * \langle \Phi_0 g_x x \rangle}; x')(x) \stackrel{(i)'}{=} \Phi_0(g_x, x).$$

$$g_x(B(\overline{f, x}; x)(x)) = \Phi_2(B(\overline{f, x * \langle B(\overline{f, x}; x)(x) \rangle}; x')) \underline{h} =$$

$$\Phi_2(B(\overline{f, x * \langle \Phi_0(g_x, x) \rangle}; x')) \underline{h} = \Phi_2(B(\overline{f, x}; x)) \underline{h}, \text{ which concludes the proof. } \square$$

We call the system of equations (1) above the **mp-system corresponding to A and $A \rightarrow B$** .

By the reasoning above we have

Proposition 3.10 For $A, B \in \mathcal{L}(\mathbf{PA})$ with $A \in \Pi_3^0, B \in \Pi_\infty^0$ one can construct functionals $\underline{\Psi} \in \widehat{\mathbf{PR}}^\omega + BR_{0,1}$ which uniformly solve the corresponding mp-system.

In a separate paper we intend to investigate in greater generality what types of unification problems can be solved by (restricted forms of) bar recursion.

The use of bar recursion in proposition 3.10 is crucial as the following proposition shows:

Proposition 3.11 Even for $A \in \Pi_3^0, B \in \Pi_0^0$ there are no functionals $\underline{\Psi} \in T$ for which it is true in \mathcal{S}^ω that they solve the corresponding mp-system uniformly.

Proof: The mp-system corresponding to A and the unique prenex normal form of $A \rightarrow B$ again is identical to the system of equations emerging from the functional interpretation of the double negation shift

$$\forall x^0 \neg \neg \exists y^0 \forall z^0 A_0(x, y, z) \rightarrow \neg \neg \forall x^0 \exists y^0 \forall z^0 A_0(x, y, z).$$

So if the mp-system would be solvable in T then this double negation shift and consequently – via negative translation – $\mathbf{PA}^\omega + \mathbf{AC}_{ar}^{0,0}$ would have a functional interpretation by functionals in T (verifiable in \mathcal{S}^ω). However it is known that all $\alpha (< \varepsilon_0)$ -recursive functions are provably recursive in $\mathbf{PA}^\omega + \mathbf{AC}_{ar}^{0,0}$ whereas the definable functions in T are α -recursive with $\alpha < \varepsilon_0$ (see e.g. [4]). \square

In contrast to this result we have

Proposition 3.12 For $A \in \Pi_2^0, B \in \Pi_\infty^0$ the corresponding mp-system has a trivial solution by substitution.

Proof: The corresponding system of equations is $x =_0 \Phi_1 f \underline{h}, \Phi_0 x =_0 f x$. Take $f := \Phi_0$ and $x := \Phi_1 f \underline{h}$. \square

4 The pointwise mp-complexity for arbitrary formulas $A, B \in \mathcal{L}(\mathbf{PA})$ and the conditions (δ) and (γ)

In the following we need a slight generalization of a result due to Schwichtenberg [24],[26]⁹ on the closure of T under the **rule** of bar recursion of type 0 (and 1):

Proposition 4.1 Let $t^2[\underline{x}^0, \underline{h}^1]$ a term of T containing at most the free variables \underline{x} of type 0 and the variables \underline{h} of type level 1. Then the functional $\lambda \underline{x}, \underline{h}, z, u, n, y. B_{0,\tau}(t[\underline{x}, \underline{h}], z, u, n, y)$ is definable in T such that \mathbf{PA}^ω (and even \mathbf{HA}^ω) proves its characterizing equations.

⁹Compare also remark 3.1 in [11] for a related result.

Proof: In [26] it is proved that for all **closed** terms t, s, r of T (of appropriate types) $\lambda n, y. B_{0,\tau} trsny$ is definable in T (formalizable in \mathbf{HA}^ω). Since there is no restriction on the type τ we can replace r, s by free variables z, u observing that $B_{0,\tau} tzuny =_\tau (B_{0,\sigma} trsny)zu$ for suitable closed λ -terms r, s (σ being a corresponding type). Moreover, Schwichtenberg's proof immediately relativizes uniformly to the case where t is allowed to contain number and function parameters yielding a primitive recursive functional (in the sense of T) in these parameters and z, u, n, y (to see this one could also use the technique of elimination of free variables from section 5 of [9]). \square

Proposition 4.2 *Let $t^{0\tau_1 \dots \tau_m \rho_1 \dots \rho_l \delta_1 \dots \delta_k}$ a (closed) term of $T_1 + BR_{0,1}$, where $\delta_1 = \dots = \delta_k = 0, \deg(\rho_1) = \dots = \deg(\rho_l) = 1, \deg(\tau_1), \dots, \deg(\tau_m) \leq 3$.*

Let $\Phi_1^{\tau_1}, \dots, \Phi_m^{\tau_m}$ be closed terms in T . Then $s := \lambda \underline{x}^{\delta} \underline{h}^{\rho}. t(\underline{x}, \underline{h}, \Phi_1, \dots, \Phi_m)$ is definable as a closed term \tilde{s} in T and $\mathbf{HA}^\omega + BR_{0,1} \vdash s =_\gamma \tilde{s}$, where γ is the type of s .

Proof: Let $t[\underline{h}]^0$ be built up from n -ary function variables \underline{h} , the combinators Π, Σ (of arbitrary finite type), $0^0, S^{00}$, closed terms $\Phi_1^{\tau_1}, \dots, \Phi_m^{\tau_m} \in T$ with $\deg(\tau_i) \leq 3$ and $B_{0,1}$. We show that $\lambda \underline{h}. t[\underline{h}]$ can be defined in T (note that this proves the proposition since the type level of R_1 is 3 and R_0 has type level 2).

For notational simplicity we assume that $\tau_i = 3$ for $i = 1, \dots, m$. By 'logical normalization' we perform all possible Π, Σ -reductions on $t[\underline{h}]^0$ and denote the result by $\hat{t}[\underline{h}]^0$ (note that $\mathbf{HA}^\omega \vdash t[\underline{h}] =_0 \hat{t}[\underline{h}]$).¹¹ The outmost constant or variable of $\hat{t}[\underline{h}]^0$ cannot be Π or Σ since if $\hat{t}[\underline{h}]^0 \equiv \Pi t_1 t_2 \dots t_i$ (resp. $\Sigma t_1 t_2 \dots t_j$) then $i \geq 2$ (resp. $j \geq 3$) since $\hat{t}[\underline{h}]$ is of type 0. But this contradicts the fact that all possible Π, Σ -reductions have been carried out already. Hence $\hat{t}[\underline{h}] \equiv 0^0, \hat{t}[\underline{h}] \equiv S(\hat{t}[\underline{h}]), \hat{t}[\underline{h}] \equiv \Phi_i^3(t_0[\underline{h}]), \hat{t}[\underline{h}] \equiv (h_i(t_1[\underline{h}]) \dots (t_j[\underline{h}]))^0$ or $\hat{t}[\underline{h}] \equiv B_{0,1}(t_1[\underline{h}]) \dots (t_6[\underline{h}])$. By proposition 4.1 (to be used in the last case only), $\hat{t}[\underline{h}]$ is primitive recursive (in the sense of T) in \underline{h} if $\hat{t}[\underline{h}]$ or $(t_0[\underline{h}]f_0)^0$ or $(t_1[\underline{h}])^0, \dots, (t_j[\underline{h}])^0$ resp. $(t_1[\underline{h}]\underline{f}_1)^0, \dots, (t_6[\underline{h}]\underline{f}_6)^0$ are primitive recursive in all of these free variables. Here \underline{f}_i are the (possibly empty) tuples of variables needed to reach the ground type 0 (note that the type levels of f_0, \underline{f}_i are ≤ 1 since all the arguments of $B_{0,1}$ and Φ_i^3 have type levels ≤ 2). We now proceed with these terms instead of $t[\underline{h}]$ (note that in the case of $t_0[\underline{h}]f_0, t_i[\underline{h}]\underline{f}_i$ we again first have to carry out all possible Π, Σ -reductions since in view of the new arguments f_0, \underline{f}_i new reductions may be possible). Eventually we end up with terms which no longer contain $B_{0,1}$ and hence are primitive recursive. So $\lambda \underline{h}. t[\underline{h}]$ is a primitive

¹⁰Here we consider the terms $\Phi_i^{\tau_i}$ as primitives, i.e. we don't carry out Π, Σ -reductions on the Π, Σ -constants occurring in these terms.

¹¹Here the notation $s[\underline{h}]$ means that s contains at most free variables from \underline{h} .

recursive functional which can be written as a closed term $\tilde{s} \in T$. To see that $\mathbf{HA}^\omega + BR_{0,1} \vdash \tilde{s}\underline{h} =_0 t[\underline{h}]$ we argue as follows: Consider a term $r \equiv B_{0,1}t_1[\underline{h}] \dots t_6[\underline{h}]$, where $t_1[\underline{h}], \dots, t_6[\underline{h}]$ do not contain $B_{0,1}$. By proposition 4.1 we can find a closed term $\tilde{r} \in T$ such that $\tilde{r}\underline{h}n\alpha$ satisfies (provably in \mathbf{HA}^ω) the instance of $BR_{0,1}$ for $t_1[\underline{h}], \dots, t_3[\underline{h}], n, \alpha$. Since $BR_{0,1}$ defines $\lambda n, \alpha. B_{0,1}(t_1[\underline{h}], \dots, t_3[\underline{h}], n, \alpha)$ uniquely in $t_i[\underline{h}]$ (provable using extensionality and bar induction or – classically – dependent choice) we have $\tilde{r}\underline{h}n\alpha =_1 B_{0,1}(t_1[\underline{h}], \dots, t_3[\underline{h}], n, \alpha)$ for all n^0, α^1 . This can be formalized in e.g. $\mathbf{PA}^\omega + \Pi_\infty^0\text{-DC}_0 + BR_{0,1}$ (where $\Pi_\infty^0\text{-DC}_0$ is the axiom schema of dependent choice of type 0 for arithmetical predicates) using the facts that all primitive recursive functionals of type 2 are \mathbf{HA}^ω -provable extensional (see [31](2.7.4)) and that $=_1 \in \Pi_1^0$ is arithmetical. Hence $\mathbf{PA}^\omega + \Pi_\infty^0\text{-DC}_0 + BR_{0,1} \vdash \tilde{r}\underline{h}n\alpha =_1 B_{0,1}(t_1[\underline{h}], \dots, t_3[\underline{h}], n, \alpha)$. But $\mathbf{PA}^\omega + \Pi_\infty^0\text{-DC}_0 + BR_{0,1} = \mathbf{PA}^\omega + \text{AC}_{ar}^{0,0} + BR_{0,1} = \mathbf{PA}^\omega + \Pi_1^0\text{-AC}^{0,0} + BR_{0,1}$ has a functional interpretation in $\mathbf{HA}^\omega + BR_{0,1}$ and hence $\mathbf{HA}^\omega + BR_{0,1} \vdash \tilde{r}\underline{h}n\alpha =_1 B_{0,1}(t_1[\underline{h}], \dots, t_3[\underline{h}], n, \alpha)$. Thus for $\hat{r}\underline{h} \equiv \tilde{r}(\underline{h}, t_4[\underline{h}], t_5[\underline{h}], t_6[\underline{h}])$ we have $\mathbf{HA}^\omega + BR_{0,1} \vdash \hat{r}\underline{h} =_0 B_{0,1}t_1[\underline{h}] \dots t_6[\underline{h}]$. The claim now follows inductively by the normalization argument above using the quantifier-free rule of extensionality of $\mathbf{HA}^\omega + BR_{0,1}$. \square

Remark 4.3 1) Proposition 4.2 is related to a result from [10] (thm.3.2 and remark 1) which in our terminology states that every term (containing only variables type of level ≤ 1) of type level ≤ 2 in $T_1 + BR_{0,1}$ has computation size strictly less than ε_0 .

2) Even for closed terms t^1 proposition 4.2 is false for $\widehat{\mathbf{PR}}^\omega + BR_{1,1}$ or $\widehat{\mathbf{PR}}^\omega + BR_{0,2}$ instead of $T_1 + BR_{0,1}$: the system $\widehat{\mathbf{PA}}^\omega \upharpoonright + \Sigma_1^1\text{-DC}$ has¹² (via negative translation) a functional interpretation in $\widehat{\mathbf{PR}}^\omega + BR_{1,1}$. But the system is proof-theoretically stronger than \mathbf{PA} (see e.g. [2] pp. 128-129) and proves more recursive functions to be total than are definable in T . The counterexample for $\widehat{\mathbf{PR}}^\omega + BR_{0,2}$ follows from the fact that $BR_{1,1}$ can be reduced to $BR_{0,2}$ (see [18],[12]). The essential formal difference between $BR_{0,1}$ and both $BR_{1,1}, BR_{0,2}$ is that the corresponding bar recursor constant $B_{0,1}$ is of type level 3 whereas both $B_{1,1}$ and $B_{0,2}$ are of type level 4 (see also [10], appendix 2).

3) Even for closed terms t^1 proposition 4.2 is false for $T_2 + BR_{0,1}$ instead of $T_1 + BR_{0,1}$. This follows from the fact that R_ρ with $\text{deg}(\rho) = 2$ (which has type level 4) can be used to iterate $B_{0,1}$ as a type-3-level functional which goes beyond $\alpha (< \varepsilon_0)$ -recursion. In fact $T_2 + BR_{0,1}$ corresponds to $\mathcal{T}_{3,4}$ in [10] where it is shown that the computation size of terms in $\mathcal{T}_{3,4}$ is $< \varepsilon_{\omega^\omega}$ and that this is optimal.

¹²Here $\Sigma_1^1\text{-DC}$ denotes the schema of dependent choice of type 1 restricted to Σ_1^1 -formulas.

Corollary 4.4 1) The same functionals of type level ≤ 2 are definable in T and in $\widehat{\mathbf{PR}}^\omega + BR_{0,1}$ (but their union $T + BR_{0,1}$ allows to define more functions).¹³

2) Let $\widehat{\mathbf{PA}}_1^\omega \upharpoonright$ be the extension of $\widehat{\mathbf{PA}}^\omega \upharpoonright$ obtained by adding the Gödel recursor R_1 for type-1-recursion with its axioms. Let $\mathcal{A} := \widehat{\mathbf{PA}}_1^\omega \upharpoonright + AC_{ar}^{0,0} + AC\text{-}qf$. If $\mathcal{A} \vdash \forall x^\rho \exists y^\tau A_0(x, y)$, where $\deg(\rho) \leq 1, \deg(\tau) \leq 2$ and $A_0(x, y)$ quantifier-free with only x, y as free variables, then one can extract a closed term $t \in T$ such that

$$\mathcal{S}^\omega \models \forall x^\delta A_0(x, tx).$$

3) Besides the usual functional interpretation (combined with negative translation) of \mathbf{PA} in T , \mathbf{PA} also has – via $\mathbf{PA} \subset \widehat{\mathbf{PA}}^\omega \upharpoonright + AC_{ar}^{0,0}$ – a functional interpretation in $\widehat{\mathbf{PR}}^\omega + BR_{0,1}$. Both functional interpretations are faithful w.r.t. the provably recursive functions of \mathbf{PA} whereas the interpretation in their union $T + BR_{0,1}$ is not.

Proof: 1) By proposition 4.2, every definable functional of type level ≤ 2 is definable in T . The other direction follows from the facts that the definable function(al)s of types 0 and 1 in T are just the $\alpha (< \varepsilon_0)$ -recursive ones, that all $\alpha (< \varepsilon_0)$ -recursive function(al)s of type level ≤ 2 are provably recursive in $\widehat{\mathbf{PA}}^\omega \upharpoonright + AC_{ar}^{0,0}$ (since the extension \mathbf{PA}^+ of \mathbf{PA} by function parameters is a subsystem of $\widehat{\mathbf{PA}}^\omega \upharpoonright + AC_{ar}^{0,0}$) and that this system has (via negative translation) a functional interpretation in $\widehat{\mathbf{PR}}^\omega + BR_{0,1}$ (see the proof of proposition 3.7). 2) From the fact that \mathcal{A} has (via negative translation) a functional interpretation in $T_1 + BR_{0,1}$ (see again the proof of proposition 3.7) and proposition 4.2 it follows that $\mathbf{HA}^\omega + BR_{0,1} \vdash \forall x^\delta A_0(x, tx)$ for some closed $t \in T$. The type structure of all continuous set-theoretical functionals \mathbf{C} from [23] (called \mathbf{S} by Scarpellini) is a model of $\mathbf{HA}^\omega + BR_{0,1}$. The conclusion now follows from the facts that $\mathbf{C}_0 = \mathcal{S}_0$ and $\mathbf{C}_1 = \mathcal{S}_1$ and that $\forall f \in \omega^\omega ([\Phi]_{\mathbf{C}} f = [\Phi]_{\mathcal{S}} f)$ for all closed terms $\Phi \in T$ of type 2. 3) follows from the proof of 2). \square

Using proposition 3.7 and proposition 4.2 we obtain that \mathbf{PA}^ω, T suffices for a **pointwise** n.c.i. of the modus ponens rule:

Proposition 4.5 Let A, B be prenex formulas in $\mathcal{L}(\mathbf{PA})$ and $(A \rightarrow B)^{pr}$ some prenex normal form of $A \rightarrow B$. Then there are functionals $\chi \in \widehat{\mathbf{PR}}^\omega + BR_{0,1}$ such that:

¹³Note that each closed term $t^2 \in \widehat{\mathbf{PR}}^\omega + BR_{0,1}$ represents a functional in \mathcal{S}^ω (so that the comparison with the type-2-functionals definable in T makes sense). This can be seen e.g. by interpreting t in the model of all continuous set-theoretical functionals \mathbf{C} from [23] since $\mathbf{C}_1 = \omega^\omega$.

If $\underline{\Phi}_A, \underline{\Phi}_{(A \rightarrow B)^{pr}} \in \mathcal{S}^\omega$ are continuous¹⁴ then also $\underline{\chi}(\underline{\Phi}_A, \underline{\Phi}_{(A \rightarrow B)^{pr}}) \in \mathcal{S}^\omega$ are continuous and if in addition

$$(1) \mathcal{S}^\omega \models (\underline{\Phi}_A \text{ n.c.i. } A) \wedge (\underline{\Phi}_{(A \rightarrow B)^{pr}} \text{ n.c.i. } (A \rightarrow B)^{pr}),$$

then

$$(2) \mathcal{S}^\omega \models \underline{\chi}(\underline{\Phi}_A, \underline{\Phi}_{(A \rightarrow B)^{pr}}) \text{ n.c.i. } B.$$

Furthermore if $\underline{\Phi}_A, \underline{\Phi}_{(A \rightarrow B)^{pr}}$ are closed terms of T then $\underline{\chi}(\underline{\Phi}_A, \underline{\Phi}_{(A \rightarrow B)^{pr}})$ can effectively be written as functionals in T .

As in proposition 3.7 this generalizes to the case where A, B contain function parameters $\underline{\alpha}, \underline{\beta}$ yielding $\underline{\chi}$ as functionals in $\underline{\Phi}_A, \underline{\Phi}_{(A \rightarrow B)^{pr}}, \underline{\alpha}, \underline{\beta}$ with $\lambda \underline{\alpha}, \underline{\beta}. \underline{\chi}(\underline{\Phi}_A, \underline{\Phi}_{(A \rightarrow B)^{pr}}, \underline{\alpha}, \underline{\beta}) \in T$ if $\underline{\Phi}_A, \underline{\Phi}_{(A \rightarrow B)^{pr}} \in T$.

Proof: The first part follows from proposition 3.7 using the fact that the extensional type structure \mathbf{C} of all continuous functionals from [23] (denoted by \mathbf{S} in [23]) is a model of $\mathbf{PA}^\omega + BR_{0,1}$, $\mathbf{C}_1 = \mathcal{S}_1$, $\mathbf{C}_2 \subset \mathcal{S}_2$ and the fact that ' $\underline{\Phi}_A, \underline{\Phi}_{(A \rightarrow B)^{pr}} \in \mathcal{S}^\omega$ continuous' iff ' $\underline{\Phi}_A, \underline{\Phi}_{(A \rightarrow B)^{pr}} \in \mathbf{C}$ ' since the type levels of these functionals are ≤ 2 .

The second part follows using proposition 4.2 \square

Remark 4.6 Note that for every \mathbb{N} -true prenex formula $A \in \mathcal{L}(\mathbf{PA})$ there are always continuous functionals $\underline{\Phi}_A \in \mathcal{S}^\omega$ satisfying the n.c.i. of A : apply unbounded search to find the least $\langle x_1, \dots, x_k \rangle$ such that $A_0(x_1, h_1 x_1, \dots)$. Furthermore by **bounded** search one can construct uniformly in $\underline{\Phi}_A$ functionals $\underline{\Phi}_A^*$ such that the implication

$$(\underline{\Phi}_A \text{ n.c.i. } A) \Rightarrow (\underline{\Phi}_A^* \text{ n.c.i. } A) \wedge (\underline{\Phi}_A^* \text{ are continuous})$$

holds for all $\underline{\Phi}_A \in \mathcal{S}^\omega$.

Proposition 4.5 implies the following result (which does not follow from the approaches to the no-counterexample interpretation via cut-elimination or functional interpretation)

Corollary 4.7 The n.c.i. of \mathbf{PA} in T (or -equivalently - by all $\alpha (< \varepsilon_0)$ -recursive functionals) satisfies Kreisel's condition (δ) in the following sense: Let $A, B \in \mathcal{L}(\mathbf{PA})$ be prenex. If

$$\underline{\Phi}_A \text{ n.c.i. } A \text{ is true for some tuple of closed terms } \underline{\Phi}_A \in T$$

and

$$\mathbf{PA} \vdash A \rightarrow B,$$

then one can construct $\underline{\Phi}_B \in T$ such that

$$\underline{\Phi}_B \text{ n.c.i. } B \text{ is true.}$$

¹⁴Here we assume that $\underline{\Phi}_A, \underline{\Phi}_{(A \rightarrow B)^{pr}}$ have the appropriate types to make them candidates for the n.c.i. of A resp. $(A \rightarrow B)^{pr}$.

Proof: By functional interpretation one extract $\underline{\Phi}_{(A \rightarrow B)^{pr}} \in T$ such that

$$\mathbf{PA}^\omega \vdash \underline{\Phi}_{(A \rightarrow B)^{pr}} \text{ n.c.i. } (A \rightarrow B)^{pr},$$

where $(A \rightarrow B)^{pr}$ is any prenex normal form of $A \rightarrow B$. The corollary now follows from 4.5. \square .

Remark 4.8 *By inspecting carefully the instances of $BR_{0,1}$ used in the proof of proposition 3.7 for given $\underline{\Phi}_A, \underline{\Phi}_{(A \rightarrow B)^{pr}} \in T$ and using the fact that proposition 4.1 can be formalized in \mathbf{HA}^ω there should be no problem to obtain corollary 4.7 also as a rule w.r.t. to \mathbf{PA}^ω -provability, i.e.*

$$\mathbf{PA}^\omega \vdash \underline{\Phi}_A \text{ n.c.i. } A \Rightarrow \mathbf{PA}^\omega \vdash \underline{\Phi}_B \text{ n.c.i. } B \text{ (even } \mathbf{HA}^\omega \vdash \underline{\Phi}_B \text{ n.c.i. } B).$$

However we will not spell out the details here.

In [17], Kreisel gives a definition of an ‘interpretation of a theory \mathcal{T} in a constructive system F ’ which essentially replaces the condition (δ) from his previous definition in [15] by the requirement

(δ') : the interpretation A_n of A implies A logically.

Kreisel mentions in his discussion in remark 2.2 of [17] that this condition is satisfied for the n.c.i. of predicate logic as well as of \mathbf{PA} only if ‘logically’ is understood in the sense of (classical) second-order logic, i.e. first-order logic extended by function quantifier **and** the axiom of choice schema

$$\forall x \exists y A(x, y) \rightarrow \exists f \forall x A(x, fx).$$

(δ) does not imply (δ') since the trivial interpretation mentioned in [15](pp.248-249) satisfies (δ) but not (δ') . In the other direction (δ') does not imply (δ) either. E.g. the n.c.i. of \mathbf{PA}_{n+1} in T_n trivially satisfies (δ') (again in the sense of second-order logic) but does not satisfy (δ) by corollary 2.4 above.

In [19] Kreisel formulates both (δ') and a version of (δ) which reads as follows ‘Having made a guess at A^* , which, in the case of the no-counterexample interpretation above is $A_0(F, f)$, we try to find, for each axiom A , a functional s_A such that

$$\mathcal{P} \vdash A^*(s_A, t) \text{ for variable } t,$$

and for each rule of inference, deriving A from \overline{A} and $\overline{\overline{A}}$ say, a functor Φ_A such that $\mathcal{P} \vdash A^*[\Phi_A(s_1, s_2), t]$ holds provided both $\mathcal{P} \vdash \overline{A}^*(s_1, t_1)$ and $\mathcal{P} \vdash \overline{\overline{A}}^*(s_2, t_2)$ hold (for variables t_1 and t_2 of appropriate type)’ (p. 378).

This is stronger than the previous formulation of (δ) from [15] by requiring the existence of **functors** which perform the rules pointwise (for provably correct interpretations for \overline{A}

and $\overline{\overline{A}}$) but also weaker by assuming the \mathcal{P} -provability of these interpretations. However the latter does not change the failure of the condition for the n.c.i. of \mathbf{PA}_{n+1} in T_n .

Finally we consider Kreisel's condition (γ) from his definition of an interpretation as given in [15]. For the n.c.i. of \mathbf{PA} in T this condition spells out as follows

$$(\gamma) \left\{ \begin{array}{l} \text{If } A \equiv \exists x_1 \forall y_1 \dots \exists x_k \forall y_k A_0(x_1, y_1, \dots, x_k, y_k, \underline{a}) \in \mathcal{L}(\mathbf{PA}) \text{ and } \mathbf{PA} \vdash \neg A. \\ \text{Then constructively it holds that for all functionals } \underline{\Phi} \in T \text{ (of suitable types)} \\ \text{there are } \underline{h} \text{ such that } A_0(\underline{\Phi}_1 \underline{h}, h_1(\underline{\Phi}_1 \underline{h}), \dots, \underline{\Phi}_k \underline{h}, h_k(\underline{\Phi}_1 \underline{h}, \dots, \underline{\Phi}_k \underline{h}), \underline{a}) \text{ is false.} \end{array} \right.$$

Similar to (δ) also the condition (γ) does not follow from the approach to the n.c.i. via cut-elimination or functional interpretation. In [16] Kreisel gave a complicated proof of (γ) using again the method of ε -substitution from [1]. We now prove a new uniform version of (γ) :

Proposition 4.9 *Let $A \equiv \exists x_1 \forall y_1 \dots \exists x_k \forall y_k A_0(x_1, y_1, \dots, x_k, y_k, \underline{a}) \in \mathcal{L}(\mathbf{PA})$ and $\mathbf{PA} \vdash \neg A$. Then from the proof of $\neg A$ one can extract functionals $\underline{\chi} \in \widehat{\mathbf{PR}}^\omega + BR_{0,1}$ such that*

$$\widehat{\mathbf{HA}}^\omega \upharpoonright + BR_{0,1} \vdash \forall \underline{\Phi} \neg A_0(\underline{\Phi}_1 \underline{h}, h_1(\underline{\Phi}_1 \underline{h}), \dots, \underline{\Phi}_k \underline{h}, h_k(\underline{\Phi}_1 \underline{h}, \dots, \underline{\Phi}_k \underline{h}), \underline{a}),$$

where $\underline{h} := \underline{\chi}(\underline{a}, \underline{\Phi})$.

Moreover for $\underline{\Phi} \in T$ (and $\widehat{\mathbf{HA}}^\omega \upharpoonright$ replaced by \mathbf{HA}^ω) one can construct a function term $t_{\underline{\Phi}} \in T$ such that $\underline{\chi}(\underline{a}, \underline{\Phi})$ can be replaced by $\tilde{\chi}(\underline{a})$, where $\tilde{\chi}_i(\underline{a}) := \lambda n. \left(\overline{\chi_i(\underline{a}, \underline{\Phi})}(t_{\underline{\Phi}}(\underline{a}, \underline{\chi}(\underline{a}, \underline{\Phi}))) \right)_n$ and $\lambda \underline{a}, n. \left(\overline{\chi_i(\underline{a}, \underline{\Phi})}(t_{\underline{\Phi}}(\underline{a}, \underline{\chi}(\underline{a}, \underline{\Phi}))) \right)_n \in T$.¹⁵

Proof: As in the proof of proposition 3.7 one has

$$\widehat{\mathbf{PA}}^\omega \upharpoonright + AC_{ar}^{0,0} \vdash \exists \underline{\Phi} \forall \underline{h} A_0(\underline{\Phi}_1 \underline{h}, h_1(\underline{\Phi}_1 \underline{h}), \dots, \underline{\Phi}_k \underline{h}, h_k(\underline{\Phi}_1 \underline{h}, \dots, \underline{\Phi}_k \underline{h}), \underline{a}) \rightarrow A.$$

Hence – using $\mathbf{PA} \vdash \neg A$ and the fact that $\mathbf{PA} \subset \widehat{\mathbf{PA}}^\omega \upharpoonright + AC_{ar}^{0,0}$ – one has

$$\widehat{\mathbf{PA}}^\omega \upharpoonright + AC_{ar}^{0,0} \vdash \forall \underline{a}, \underline{\Phi} \exists \underline{h} \neg A_0(\underline{\Phi}_1 \underline{h}, h_1(\underline{\Phi}_1 \underline{h}), \dots, \underline{\Phi}_k \underline{h}, h_k(\underline{\Phi}_1 \underline{h}, \dots, \underline{\Phi}_k \underline{h}), \underline{a}).$$

The first part of the proposition now follows from the fact that $\widehat{\mathbf{PA}}^\omega \upharpoonright + AC_{ar}^{0,0}$ has (via negative translation) a functional interpretation in $\widehat{\mathbf{HA}}^\omega \upharpoonright + BR_{0,1}$.

¹⁵Here we refer to a standard coding $\overline{fx} := \langle f0, \dots, f(x-1) \rangle$ of finite sequences.

For the second part we use that for $\underline{\Phi} \in T$ there is a function term $t_A \in T$ such that

$$t_{A\underline{a}\underline{h}} =_0 0 \leftrightarrow \neg A_0(\Phi_1\underline{h}, h_1(\Phi_1\underline{h}), \dots, \Phi_k\underline{h}, h_k(\Phi_1\underline{h}, \dots, \Phi_k\underline{h}), \underline{a}).$$

By [31](2.7.8) t_A is continuous in \underline{h} , provably in \mathbf{HA}^ω with a modulus of pointwise continuity $t \in T$, i.e.

$$\forall \underline{h}, \tilde{\underline{h}} (\bigwedge_i (\overline{h}_i(t(\underline{a}, \underline{h})) =_0 \overline{h}_i(t(\underline{a}, \underline{h}))) \rightarrow t_A(\underline{a}, \underline{h}) =_0 t_A(\underline{a}, \tilde{\underline{h}})).$$

The last claim in the proposition follows from 4.2. \square .

Remark 4.10 *Classically one can easily find functions \underline{h} satisfying (γ) for $\underline{\Phi} \in T$ (or for arbitrary continuous $\underline{\Phi}$) by unbounded recursive search:*

*Let $\underline{\Phi} \in T$. Using the continuity argument from the proof above $\exists \underline{h} \neg A_0(\dots \underline{h} \dots)$ (which **classically** follows trivially from the assumptions of (γ) as in the proof above) implies $\exists \underline{x}^0 \neg A_0(\dots \lambda n.(x_1)_n, \dots, \lambda n.(x_k)_n \dots)$. By unbounded search (recursively in \underline{a}) one finds \underline{x} realizing $\exists \underline{x}$.*

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