

# Simplicial Sets within Cubical Sets

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# Computational Meaning of HoTT? (1)

**HoTT** is **Intensional Type Theory** together with Voevodsky's **Univalence Axiom (UA)** and **Higher Inductive Types (HITs)**.

Let us recall what *UA* says. This requires a few definitions

$$\text{iscontr}(X : \text{Set}) = (\Sigma x : X)(\Pi y : X) \text{Id}_X(x, y)$$

$$\begin{aligned} \text{hfiber}(X, Y : \text{Set})(f : X \rightarrow Y)(y : Y) &= \\ &= (\Sigma x : X) \text{Id}_Y(f(x), y) \end{aligned}$$

$$\begin{aligned} \text{isweq}(X, Y : \text{Set})(f : X \rightarrow Y) &= \\ &= (\Pi y : Y) \text{iscontr}(\text{hfiber}(X, Y, f, y)) \end{aligned}$$

$$\text{Weq}(X, Y : \text{Set}) = (\Sigma f : X \rightarrow Y) \text{isweq}(X, Y, f)$$

## Computational Meaning of HoTT? (2)

Using the eliminator  $J$  for identity types one easily constructs a map

$$\text{eqweq}(X, Y : \text{Set}) : \text{Id}_{\text{Set}}(X, Y) \rightarrow \text{Weq}(X, Y)$$

Then the **Univalence Axiom**

$$\text{UA} : (\prod X, Y : \text{Set}) \text{isweq}(\text{eqweq}(X, Y))$$

postulates that all maps  $\text{eqweq}(X, Y)$  are weak equivalences.

It has been shown that **simplicial sets** provide a model of HoTT interpreting types as Kan complexes.

But UA as it is **lacks computational meaning**:  
**what should be rewrite rules for the constant UA?**

T. Coquand et.al. [CCHM] have developed a *Cubical Type Theory* with computational meaning and in which one can **derive** UA.

# From Simplicial Sets to Cubical Sets (1)

Bezem and Coquand have shown that the theory of simplicial sets is **not constructive**, e.g. one cannot show constructively that Kan complexes are closed under exponentiation!

This limitation, however, can be overcome when working in **cubical sets**.

The crucial property is that representable objects are closed under finite products. In **sSet** the interval  $\mathbb{I}$  is representable but  $\mathbb{I} \times \mathbb{I}$  is not!

The site of **sSet** is  $\Delta$ , the full subcategory of **Poset** on finite non-empty linear posets.

The site of **cSet** is  $\square$ , the full subcategory of **Poset** of finite powers of  $\mathbf{2}$ .

Splitting idempotents in  $\square$  gives rise to **FL**, the full subcategory of **Poset** on finite lattices.

## From Simplicial Sets to Cubical Sets (2)

Notice that  $\square$  is op-equivalent to the category of *free finitely generated distributive lattices*. Coquand et.al. – just for convenience – used the opposite of *free finitely generated de Morgan algebras*.

Let  $i : \Delta \rightarrow \mathbf{FL}$  be the inclusion functor. The restriction functor  $i^*$  from  $\mathbf{cSet} = \widehat{\mathbf{FL}}$  to  $\mathbf{sSet} = \widehat{\Delta}$  has left and right adjoints  $i_!$  and  $i_*$ , respectively. Since the restriction of the nerve functor  $\text{Nv}$  to  $\mathbf{FL}$  is given by  $i^* \circ Y_{\mathbf{FL}}$  we have

$$i_*(X)(L) \cong \mathbf{cSet}(Y(L), i_*(X)) \cong \mathbf{sSet}(i^*Y(L), X) \cong \mathbf{sSet}(\text{Nv}(L), X)$$

from which it follows that  $i_*$  is full and faithful (and thus also  $i_!$ ). Since  $i^*$  has a left adjoint  $i_!$  (given by left Kan extension of  $Y_{\mathbf{FL}} \circ i$  along  $Y_{\Delta}$ ) it preserves (finite) limits and thus  $i^* \dashv i_*$  is an **injective geometric morphism**.

# The topology on $\mathbf{FL}$ inducing $\mathbf{sSet}$

A sieve  $S \subseteq Y_{\mathbf{FL}}(L)$  **covers**  $L$  iff  $i^*S = i^*Y_{\mathbf{FL}}(L) = \text{Nv}(L)$  iff  $S$  contains all chains in  $L$ , i.e. all monotone maps  $[n] \rightarrow L$ .

The corresponding **closure operator**  $j : \Omega \rightarrow \Omega$  sends  $S \subseteq Y_{\mathbf{FL}}(L)$  to all  $u : K \rightarrow L$  such that  $uc \in S$  for all chains  $c : [n] \rightarrow K$ .

# Cisinski Model Structures on $\mathbf{cSet}$ and $\mathbf{sSet}$ (1)

For Cisinski model structures on  $\widehat{\mathcal{C}}$  its class of **cofibrations** consist of all monos.

Its class of trivial fibrations consists of maps weakly right orthogonal to all monos.

A **naive fibration** is a map weakly right orthogonal to **all cylinders**, i.e. monos of the form

$$(\{\varepsilon\} \times X) \cup (\mathbb{I} \times Y) \hookrightarrow \mathbb{I} \times X$$

where  $Y \subseteq X$ ,  $\varepsilon \in \{0, 1\} \subseteq \mathbb{I} = Y(2)$ .

The **trivial cofibrations** are the maps weakly left orthogonal to all naive fibrations.

On  $\mathbf{sSet}$  this construction gives the classical model structure and on  $\mathbf{cSet}$  the one employed by Coquand et.al. in [CCHM].

## Cisinski Model Structures on $\mathbf{cSet}$ and $\mathbf{sSet}$ (2)

Since  $i^*$  preserves  $\mathbb{I}$  and finite limits and it retracts monos of  $\mathbf{cSet}$  onto monos of  $\mathbf{sSet}$  it follows that  $i^*$  retracts cylinders in  $\mathbf{sSet}$  onto cylinders in  $\mathbf{sSet}$ .

Thus  $i_*p$  is a fibration in  $\mathbf{cSet}$  iff  $p$  is a fibration in  $\mathbf{sSet}$ .

Accordingly, we have  $\mathcal{F}_{\mathbf{sSet}} = \mathbf{sSet} \cap \mathcal{F}_{\mathbf{cSet}}$ .

Thus  $i^* \dashv i_*$  is a *Quillen pair*.

We don't know whether it is a Quillen equivalence.

But presumably not!



## $i_*$ preserves and reflects “equality” of fibrant objects

**Theorem** If  $A$  and  $B$  are fibrant in  $\mathbf{sSet}$ , i.e. Kan complexes, a map  $f : A \rightarrow B$  is a weak equivalence in  $\mathbf{sSet}$  iff it is a weak equivalence in  $\mathbf{cSet}$ .

*Proof.*

One can show that  $i_*(\text{hfiber}(f)) \simeq \text{hfiber}(i_*f)$  and thus

$\forall m \in \text{Mono}(\mathbf{cSet})(m \perp \text{hfiber}(i_*f))$  iff

$\forall m \in \text{Mono}(\mathbf{cSet})(m \perp i_*(\text{hfiber}(f)))$  iff

$\forall m \in \text{Mono}(\mathbf{cSet})(i^*m \perp \text{hfiber}(f))$  iff<sup>1</sup>

$\forall m \in \text{Mono}(\mathbf{sSet})(m \perp \text{hfiber}(f))$

i.e.<sup>2</sup>  $i_*f$  is a weak equivalence in  $\mathbf{cSet}$  iff  $f$  is a weak equivalence in  $\mathbf{sSet}$ . □

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<sup>1</sup>since the monos in  $\mathbf{sSet}$  are precisely the sheafifications of monos in  $\mathbf{cSet}$

<sup>2</sup>as shown by Voevodsky for fibrant objects  $A$  and  $B$  a map  $w : A \rightarrow B$  is a weak equivalence iff  $\text{hfiber}(w)$  is a trivial cofibrations, i.e.  $m \perp \text{hfiber}(w)$  for all monos  $m$

# Universes in **cSet** and **sSet**

In both **cSet** and **sSet** we may construct universes

$$\pi_c : E_c \rightarrow U_c \quad \text{and} \quad \pi_s : E_s \rightarrow U_s$$

à la Yoneda where  $U(I)$  consists of all small fibrations over  $Y(I)$ .  
In both cases one could show the universe to be fibrant and univalent!

However, in case of **sSet** this requires heavy choice (due to use of minimal fibrations!) but not so for **cSet**.

If  $i^*$  preserved fibrations, i.e.  $i_!$  preserved cylinders, then  $i^*\pi_c$  were a universe generic for small Kan fibrations.

Composing this with the map sending a small fibration over  $Y_{\square}([n])$  to its sheafification which is a small fibration over  $Y_{\Delta}([n])$  we would obtain a univalent universe equivalent to  $\pi_s$ .

# But $i_1 : \mathbf{sSet} \hookrightarrow \mathbf{cSet}$ does not preserve finite limits!

Since both  $i_1$  and  $i^*$  are cocontinuous and  $i_1$  is full and faithful the full embedding  $i_1$  preserves and reflects colimits. Thus  $\mathbf{sSet}$  is a full subcategory of  $\mathbf{cSet}$  closed under colimits (taken in  $\mathbf{cSet}$ ).

By a well known theorem  $i_1$  preserves finite limits iff  $\mathbf{FL}(L, i(-))$  is **flat** for all  $L \in \mathbf{FL}$  which, however, is not the case:

let  $L = [1] \times [1]$  then there are (precisely two) 1-1 maps  $f, g : [1] \times [1] \rightarrow [3]$  not fitting into a diagram of the form

$$\begin{array}{ccccc} & & [n] & & \\ & \xleftarrow{p} & & \xrightarrow{q} & \\ [3] & & & & [3] \\ & \nwarrow f & \uparrow h & \nearrow g & \\ & & [1] \times [1] & & \end{array}$$

## Staying within **cSet** avoids the problem!

Cut  $U_c$  down to the subuniverse  $U_{cs}$  consisting at stage  $L \in \mathbf{FL}$  of all  $A \in U(L)$  such that  $a : \text{El}(A) \rightarrow Y(L)$  is a family of sheaves, i.e.

$$\begin{array}{ccc} \text{El}(A) & \xrightarrow{\eta_{\text{El}(A)}} & i_* i^* \text{El}(A) \\ \downarrow a & \lrcorner & \downarrow i_* i^* a \\ Y(L) & \xrightarrow{\eta_{Y(L)}} & i_* i^* Y(L) \end{array}$$

or equivalently  $a$  is strictly right orthogonal to all subobjects of representables which are dense, i.e. inverted by  $i^*$ .

This universe  $U_{cs}$  classifies small families of sheaves which are fibrations in **cSet**. If the indexing object is a sheaf such families coincide with images of fibrations under  $i_*$ .

# Conclusion and Outlook

- We have shown that simplicial sets form a(n essential) subtopos of cubical sets.
- Moreover, **sSet** is a submodel of **cSet** since the inclusion  $i_* : \mathbf{sSet} \hookrightarrow \mathbf{cSet}$  preserves  $\Sigma$  and  $\Pi$  and also the interval  $\mathbb{I}$  and thus also identity types.
- Cutting down the universe  $U_c$  to families of sheaves weakly classifies (small) fibrations of sheaves.